Shared Counters and Parallelism



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Based on slides by Maurice Herlihy

and Nir Shavit

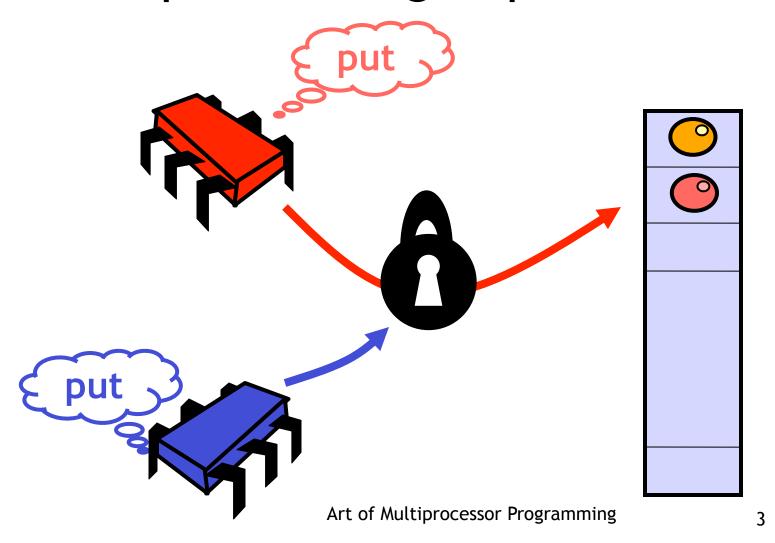
A Shared Pool

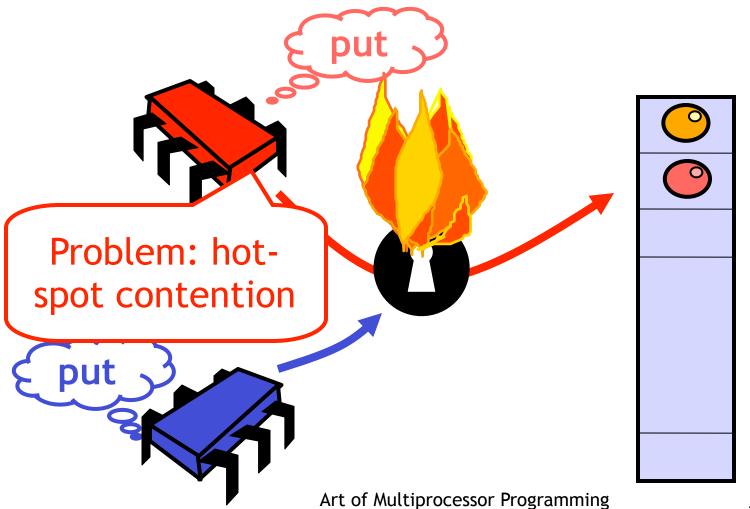
```
public interface Pool {
  public void put(Object x);
  public Object remove();
}
```

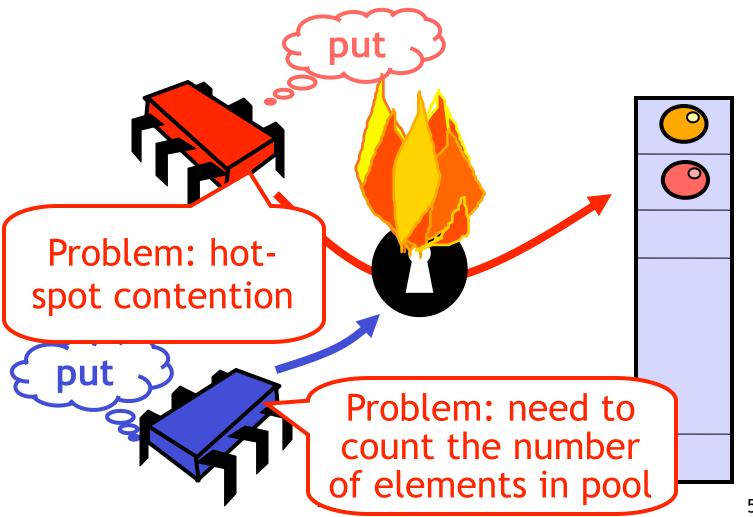
Unordered set of objects

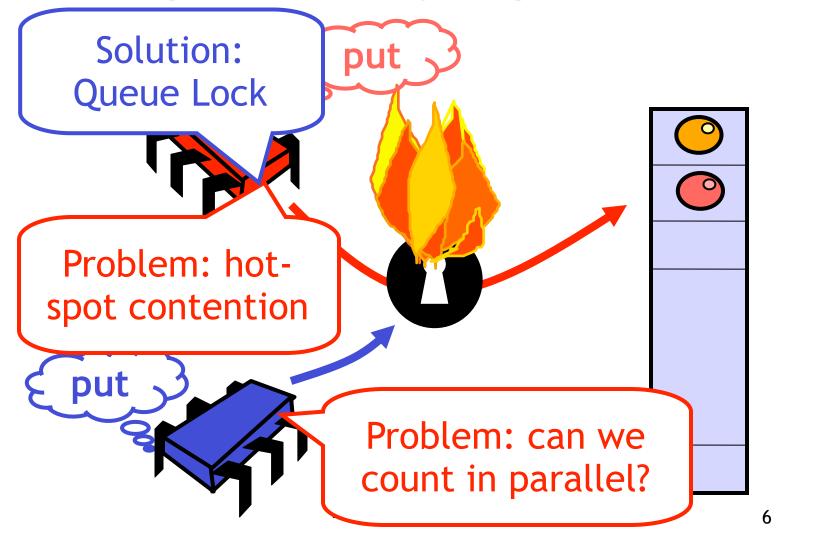
- Put
 - Inserts object
 - blocks if full

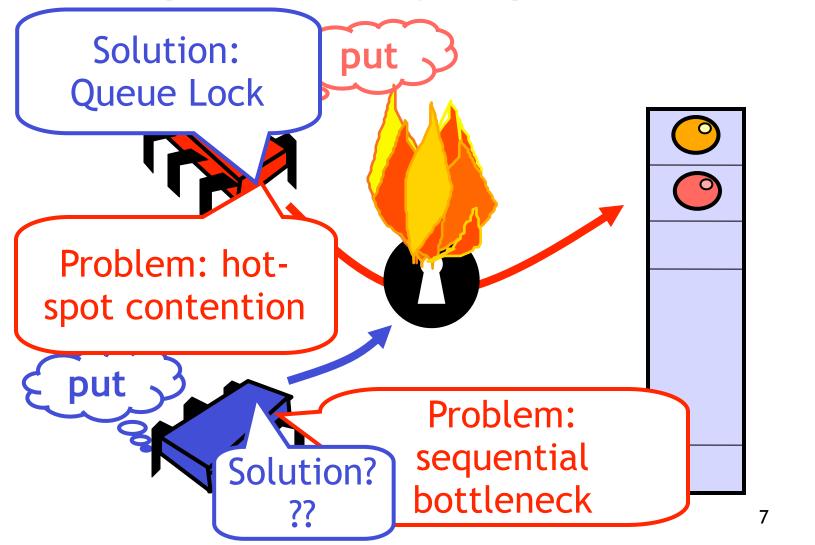
- Remove
 - Removes & returns an object
 - blocks if empty



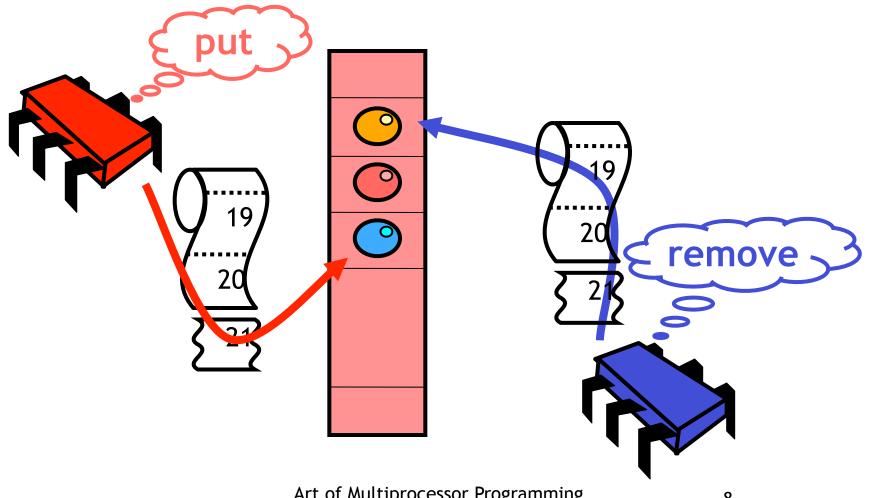




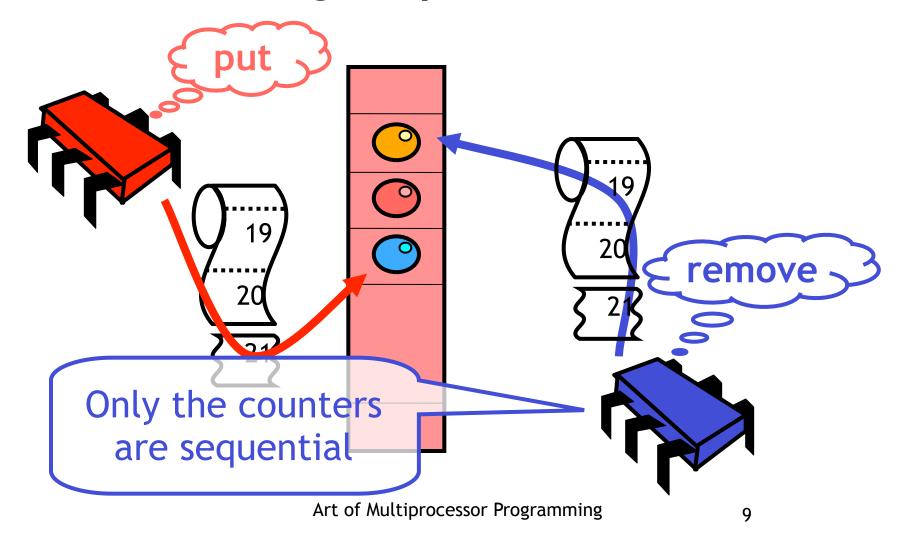


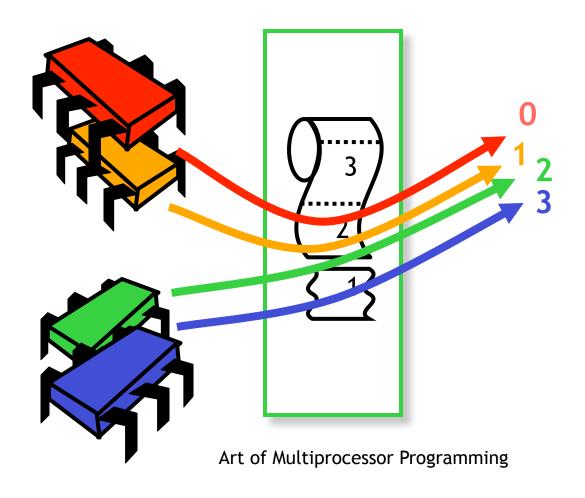


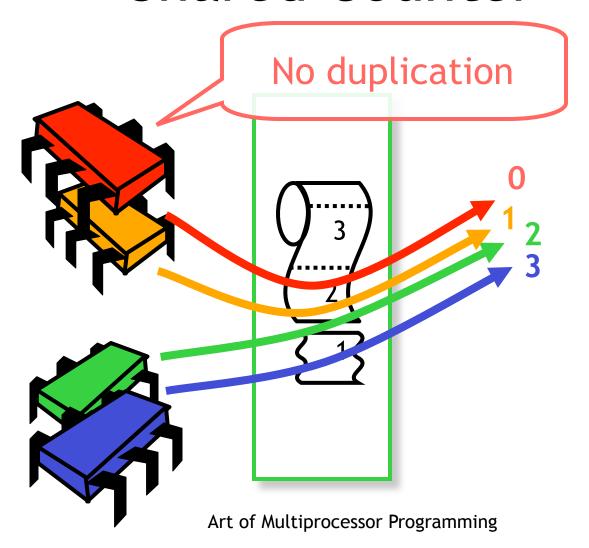
Counting Implementation

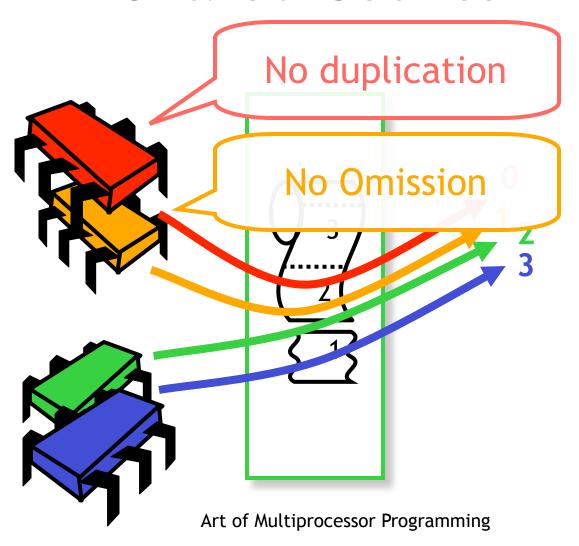


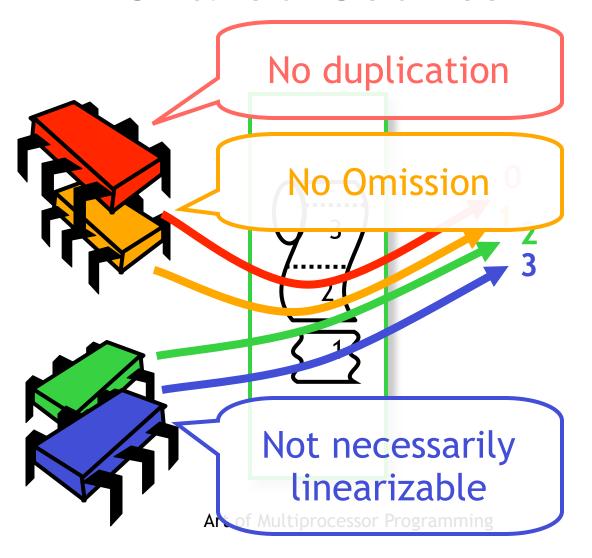
Counting Implementation



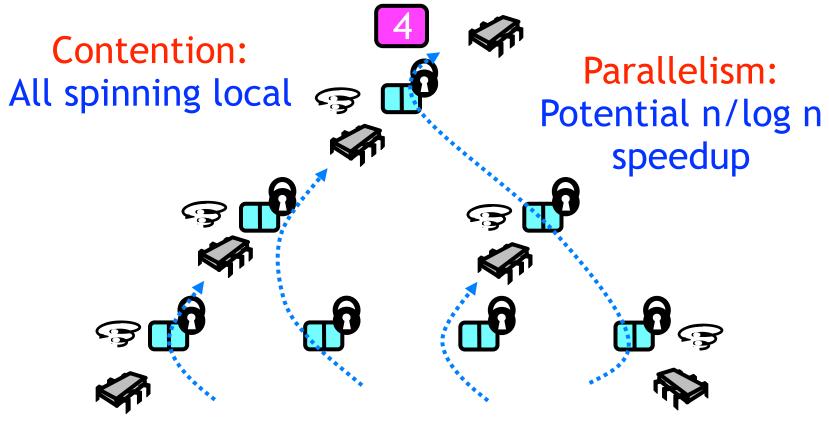


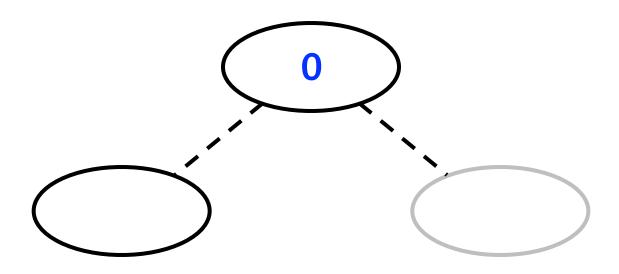


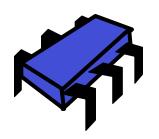


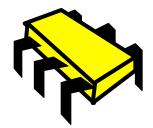


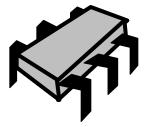
- Can we build a shared counter with
 - Low memory contention, and
 - Real parallelism?
- Locking
 - Can use queue locks to reduce contention
 - No help with parallelism issue ...

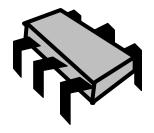


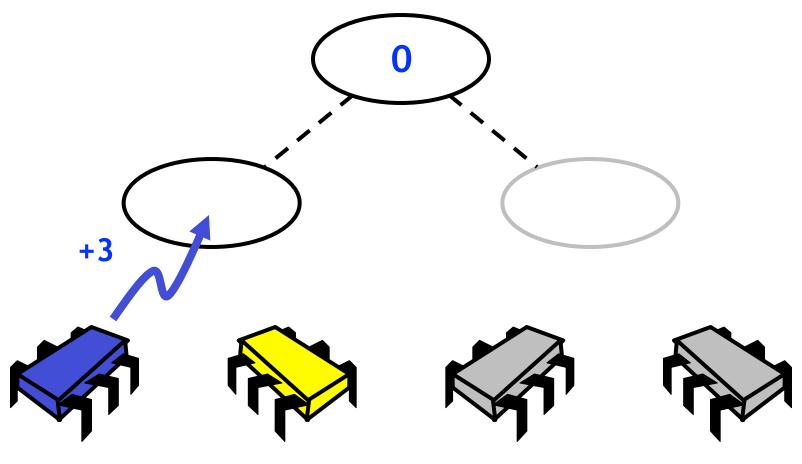


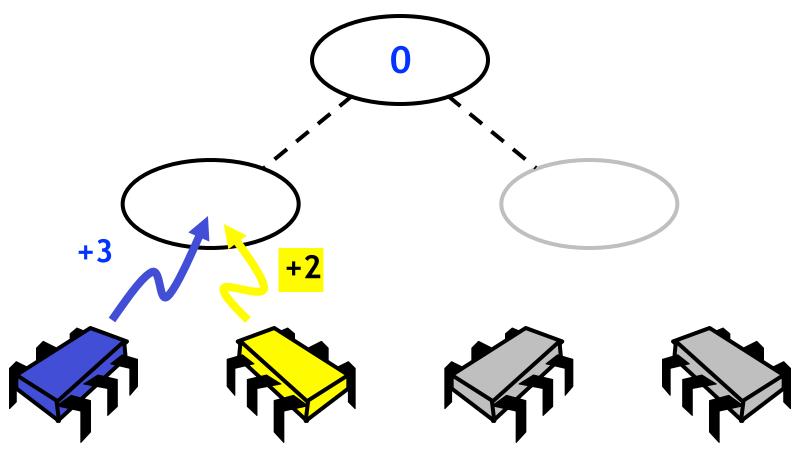


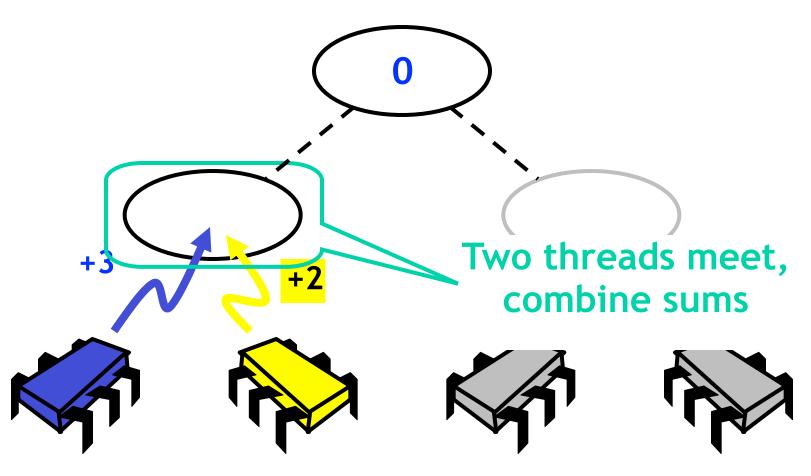


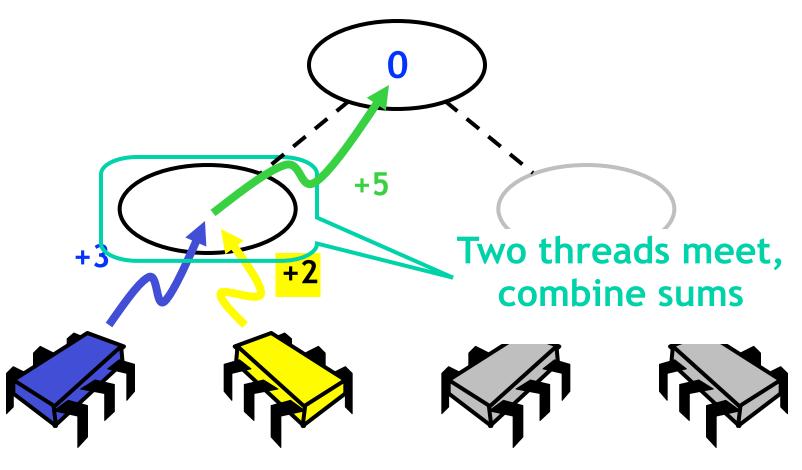


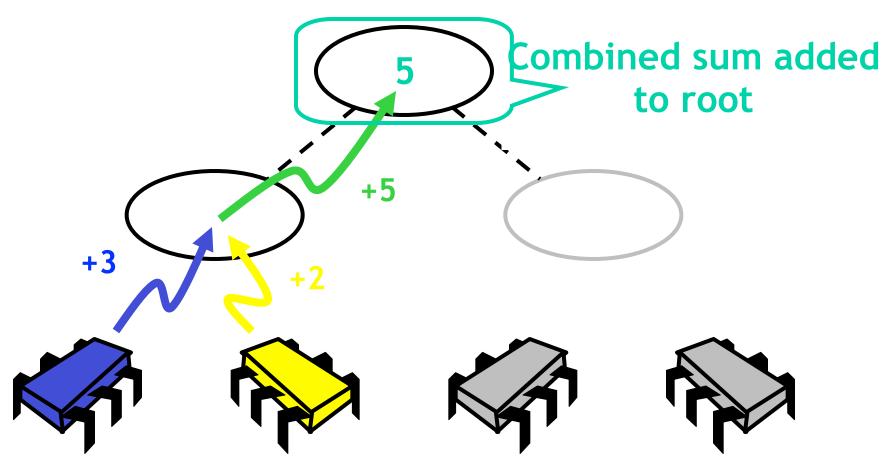


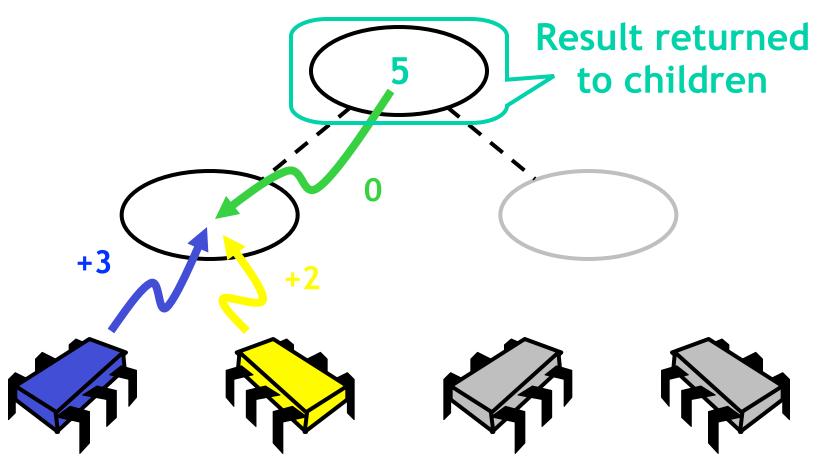


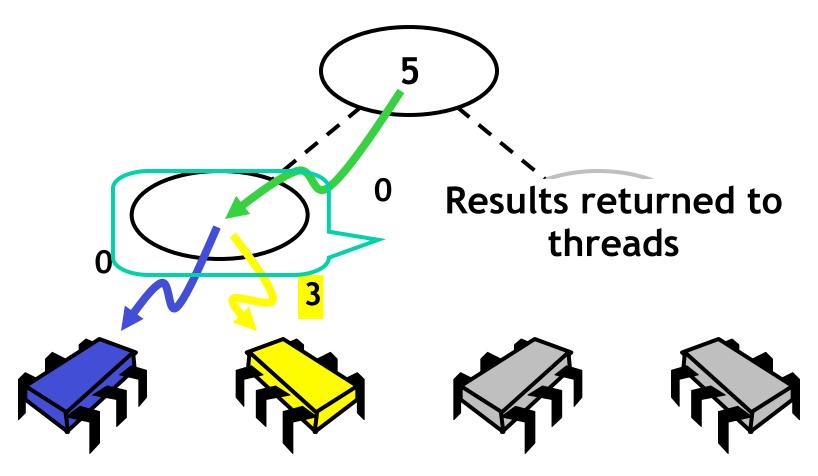






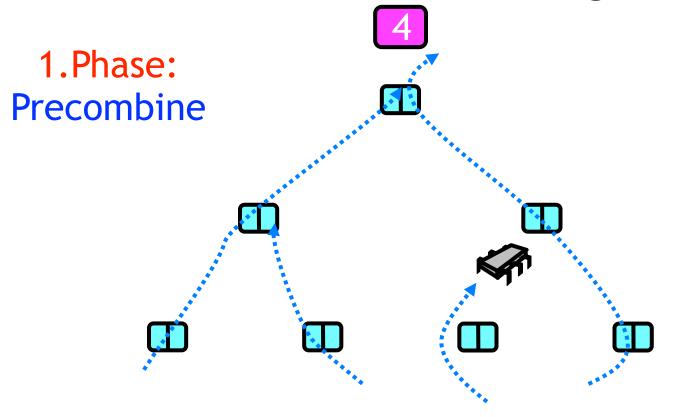


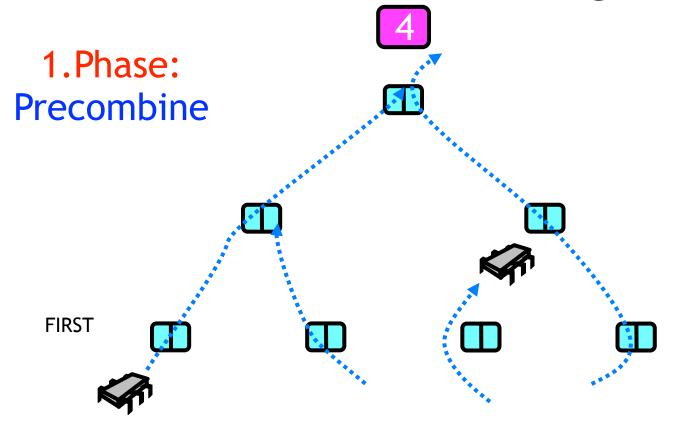


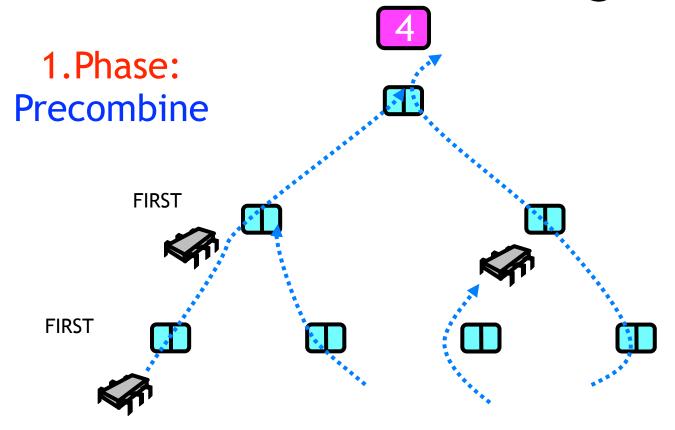


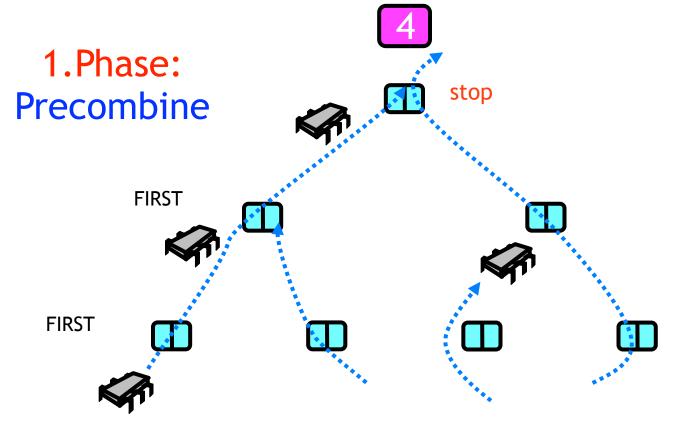
Devil in the Details

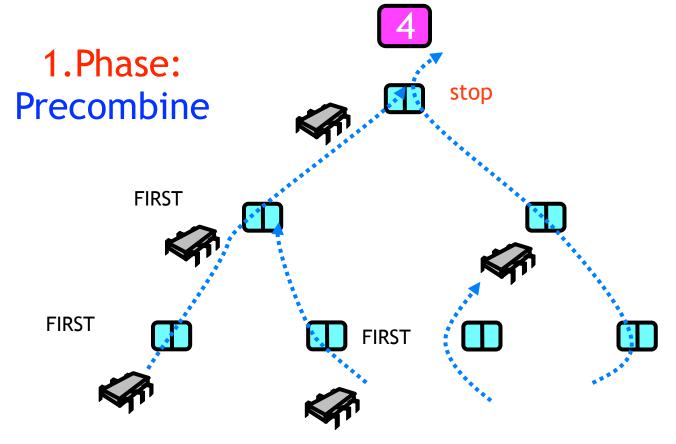
- What if
 - threads don't arrive at the same time?
- Wait for a partner to show up?
 - How long to wait?
 - Waiting times add up ...
- Instead
 - Use multi-phase algorithm
 - Try to wait in parallel ...

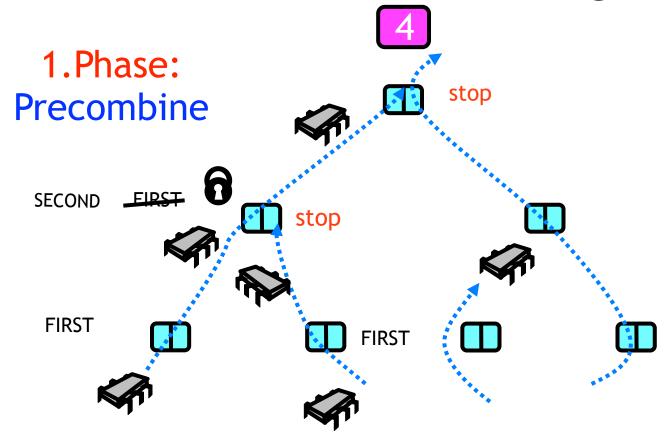






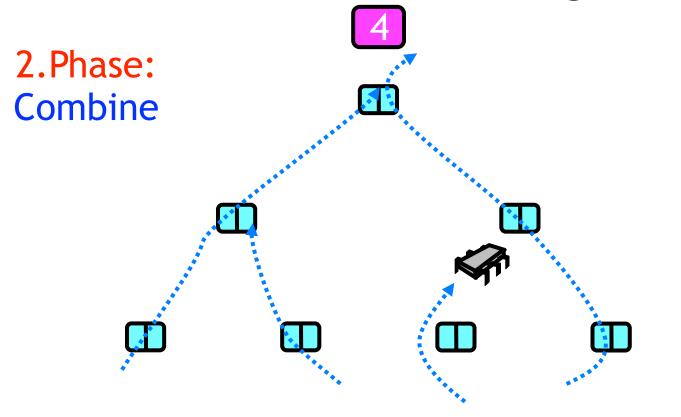


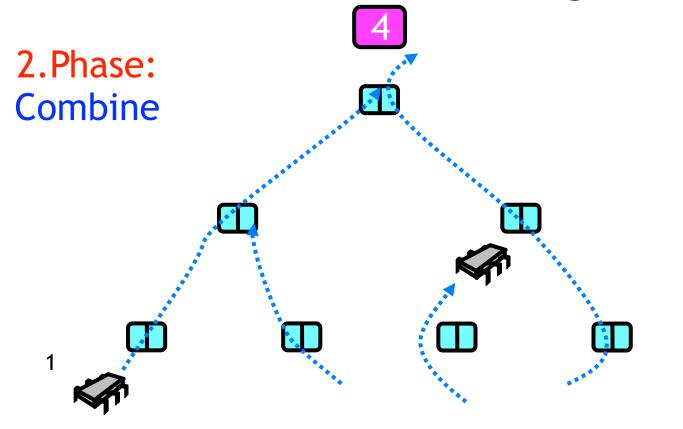


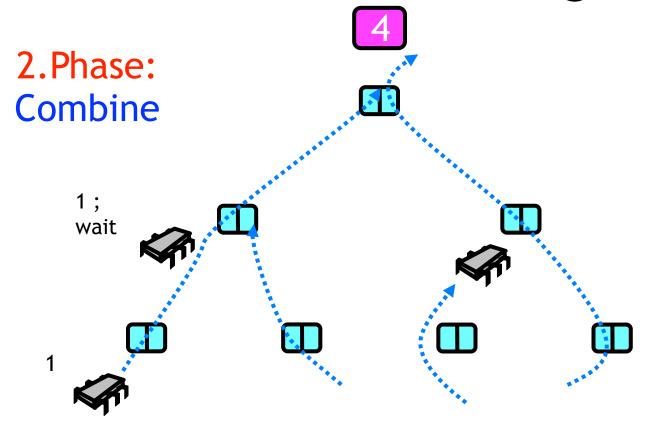


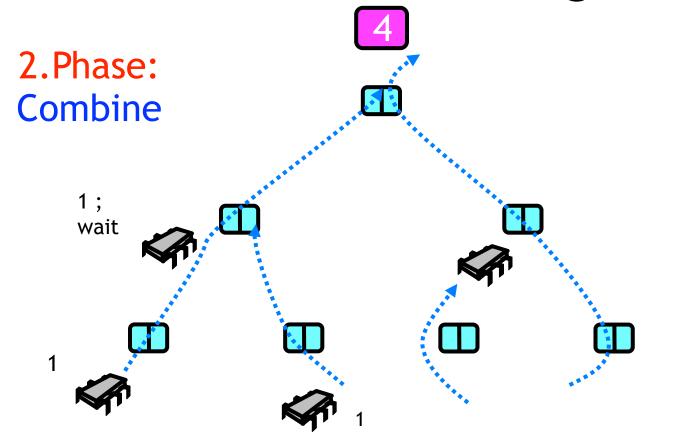
Phases

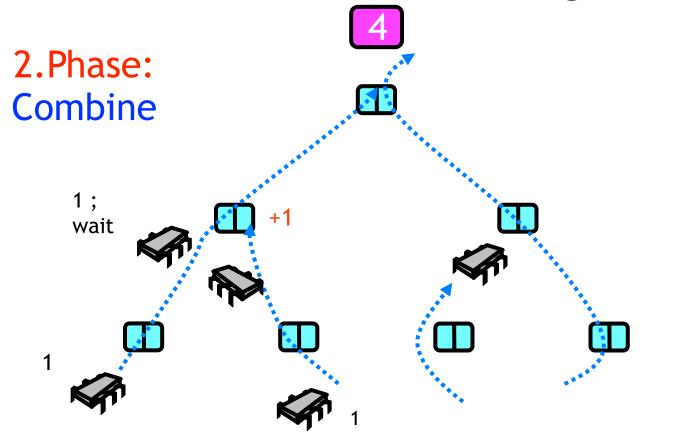
- Precombining
 - Set up combining rendezvous
- Combining
 - Collect and combine operations
- Operation
 - Hand off to higher thread
- Distribution
 - Distribute results to waiting threads

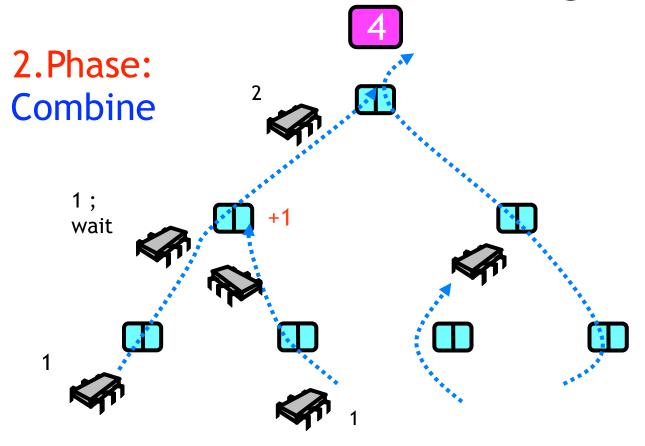


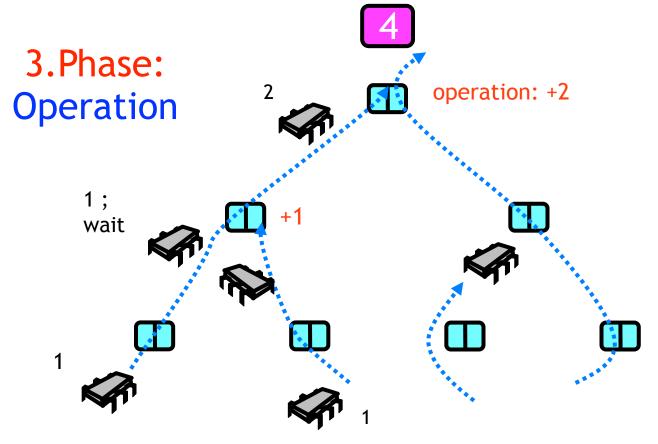


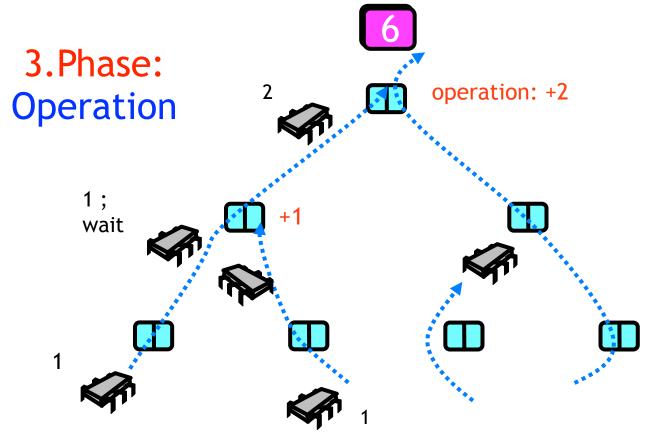


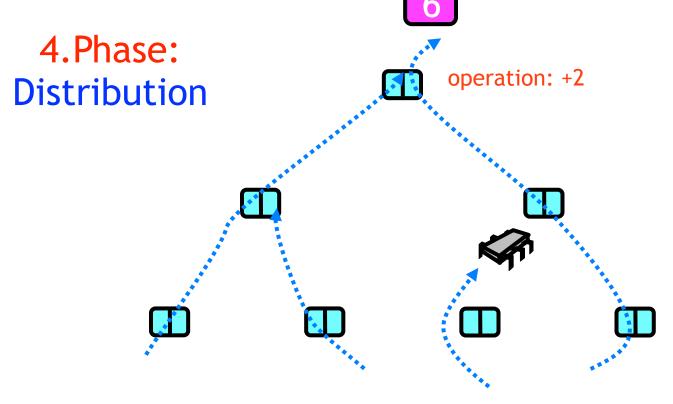


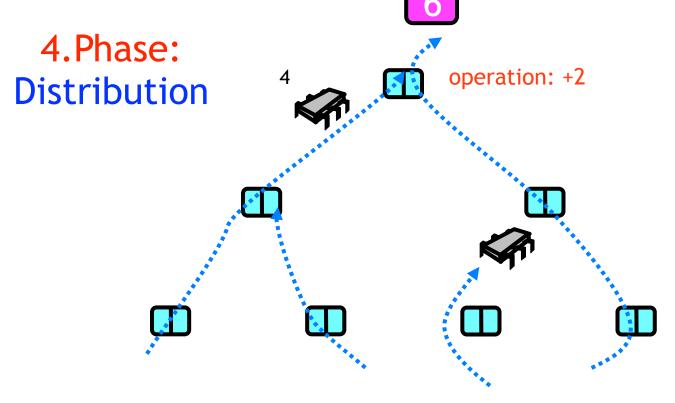


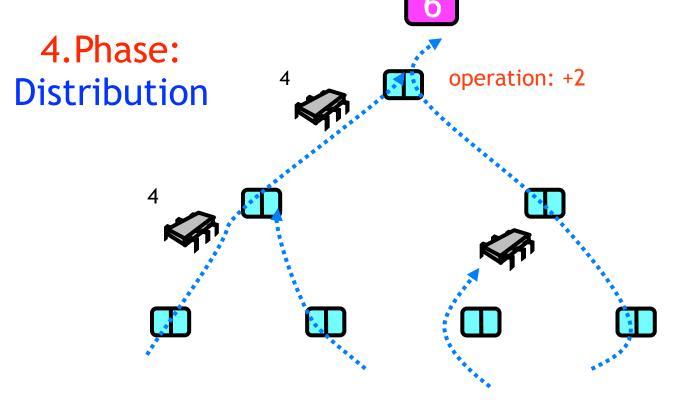


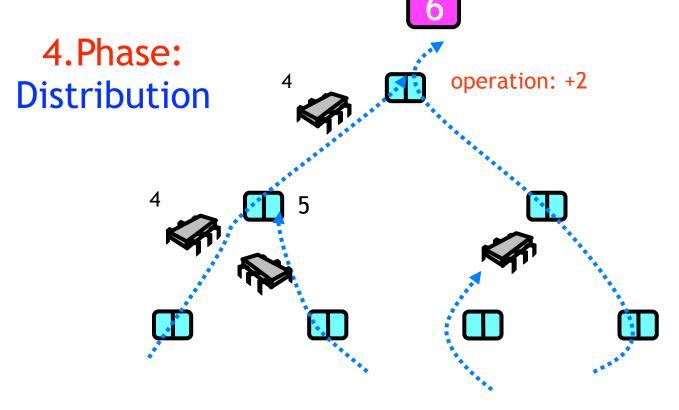


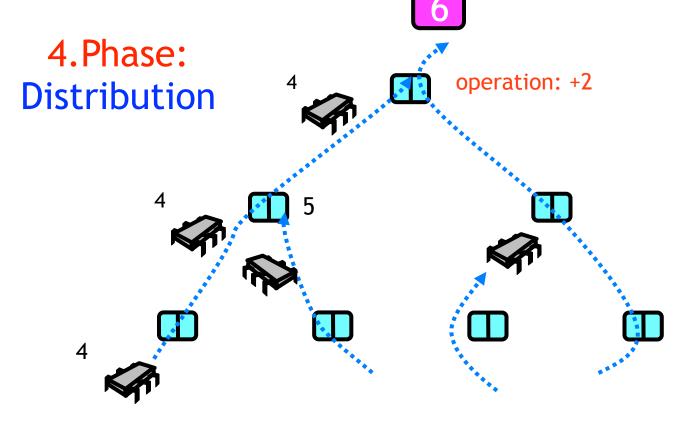


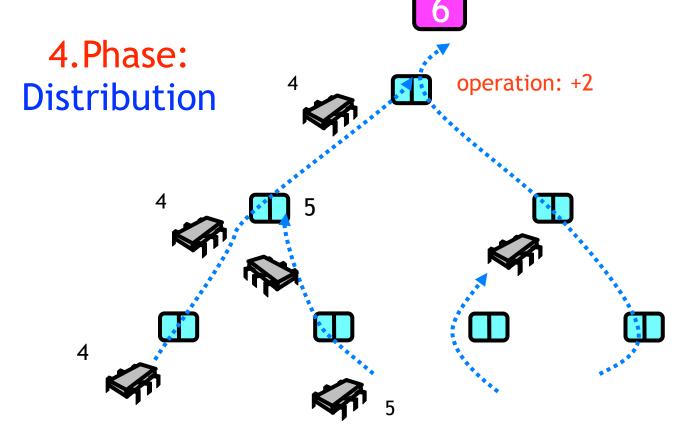












```
enum CStatus{
IDLE, FIRST, SECOND, DONE, ROOT};
```

```
enum CStatus{

IDLE, FIRST, SECOND, DONE, ROOT};

Nothing going on
```

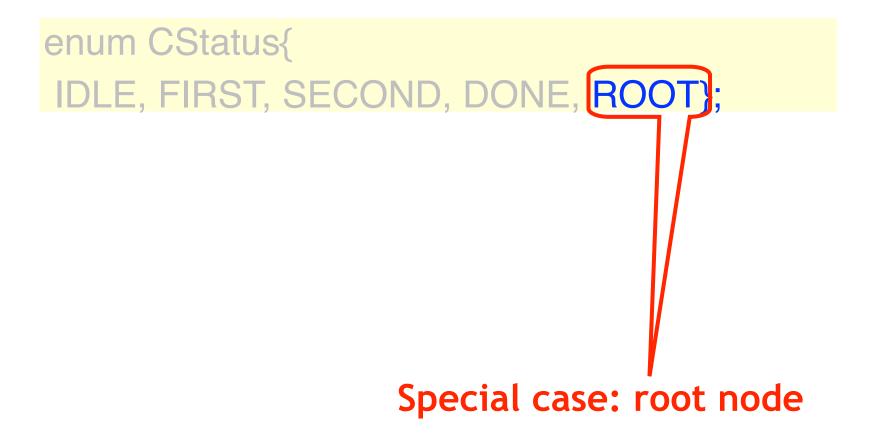


1st thread is partner for combining, will return soon to check for 2nd thread



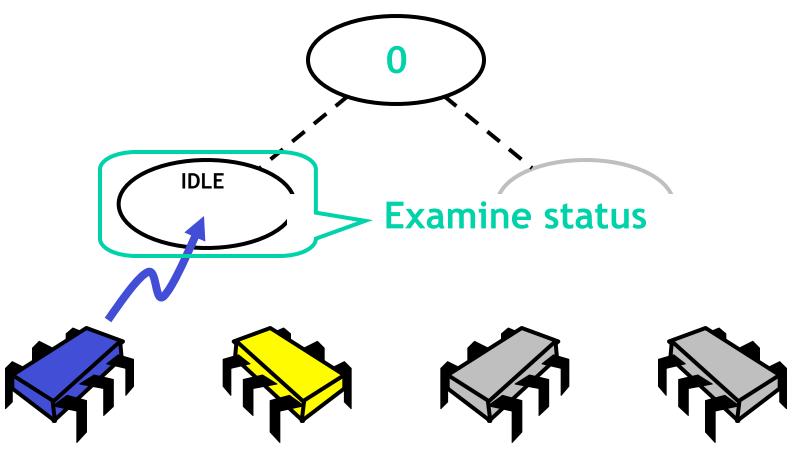


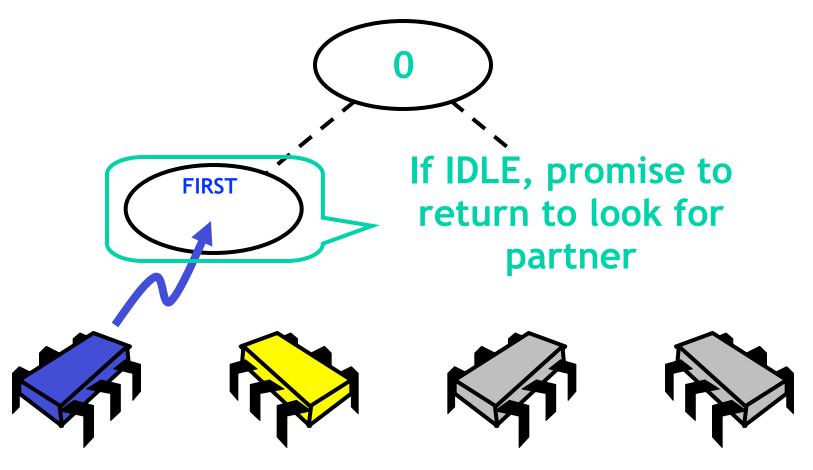
1st thread has completed operation & deposited result for 2nd thread

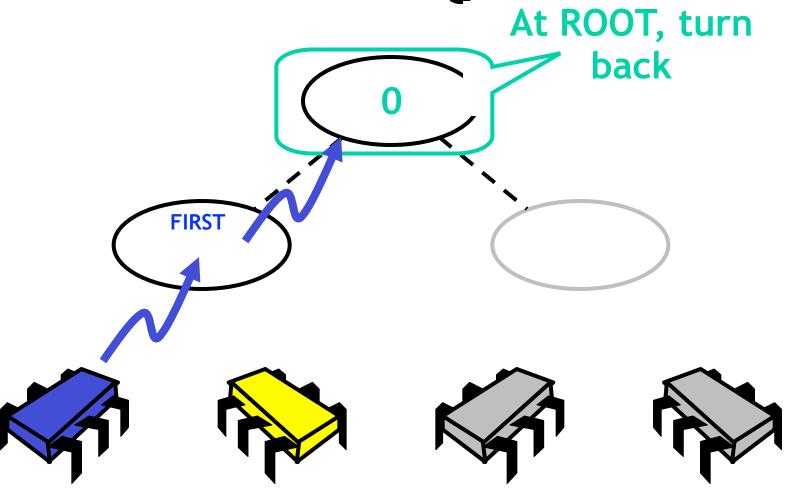


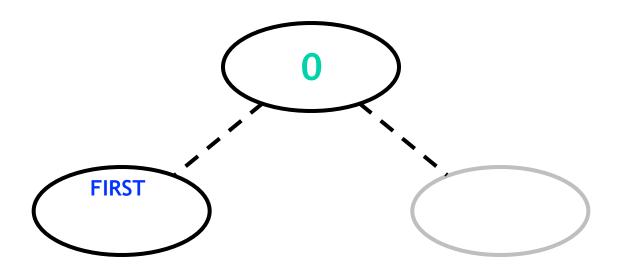
Node Synchronization

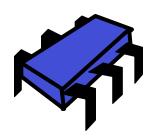
- Short-term
 - Synchronized methods
 - Consistency during method call
- Long-term
 - Boolean locked field
 - Consistency across calls

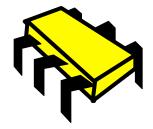


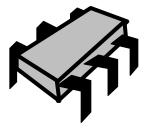


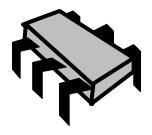


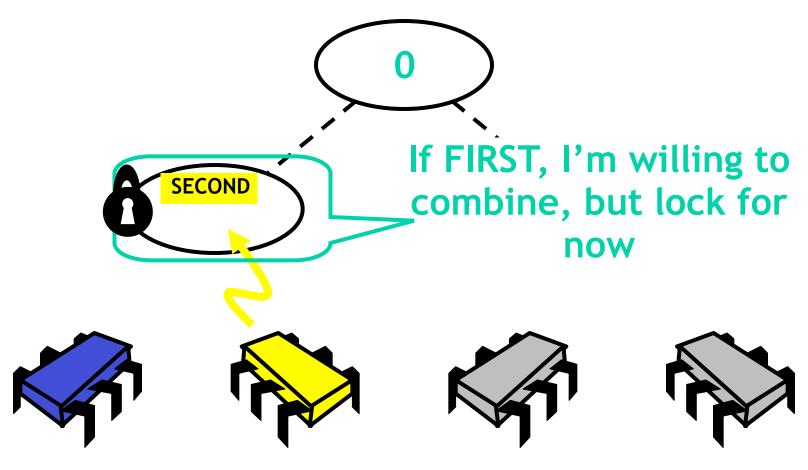












Code

- Tree class
 - In charge of navigation
- Node class
 - Combining state
 - Synchronization state
 - Bookkeeping

```
Node node = myLeaf;
while (node.precombine()) {
  node = node.parent;
  }
Node stop = node;
```

```
Node node = myLeaf;
while (node.precombine)
 node = node.parer
Node stop = node;
                              Start at leaf
```

```
Node node = myLeaf;
while (node.precombine()) {
 node = node.parent;
Move up while instructed
```

```
Node node = myLeaf;
while (node.precombine()) {
  node = node.parent;
  }
Node stop = node;
```

Remember where we stopped

Precombining Node

```
synchronized boolean precombine() {
while (locked) wait();
switch (cStatus) {
 case IDLE: cStatus = CStatus.FIRST;
        return true;
 case FIRST: locked = true;
        cStatus = CStatus.SECOND;
        return false;
 case ROOT: return false;
 default: throw new PanicException()
```

Precombining Node

```
synchronized boolean precombine() {
while (locked) wait();
switch (cStatus
 case IDLE: cStatus = CStatus.FIRST;
       return true;
 case FIRST: locked =
        cStatus = CStatus
        return false;
 case ROOT: return false;
 default: throw new PanicException()
                                   Short-term
                                synchronization
```

Synchronization

```
synchronized boolean precombine() {
while (locked) wait();
switch (cStatus) {
case IDLE: cStatus = CStatus.FIRST;
       return true;
 case FIRST: locked = tive;
        cStatus = CStatus SE
        return false;
 case ROOT: return false;
 default: throw new PanicException()
                        Wait while node is locked
```

Precombining Node

```
synchronized boolean precombine() {
while (locked) wait():
switch (cStatus) {
 case IDLE: cStatus = CStatus.FIRST;
       return true;
 case FIRST: locked = true
        cStatus = CStatus.SECOND;
        return false:
 case ROOT: return false;
 default: throw new PanicException(
                            Check combining status
```

Node was IDLE

```
synchronized boolean precombine() {
while (locked) {wait();}
switch (cStatus) {
 case IDLE: cStatus = CStatus.FIRST;
       return true;
 case FIRST: locked = true;
        cStatus = CStatus.SECO
        return false;
 case ROOT: return false;
 default: throw new PanicException()
  I will return to look for
      combining value
```

Precombining Node

```
synchronized boolean precombine() {
while (locked) {wait();}
switch (cStatus) {
 case IDLE: cStatus = CStatus.FIRST;
       return true;
 case FIRST. locked = true;
        cStatus = CStatus.SECOND;
        return false:
                       Continue up the tree
 case ROOT: return
 default: throw new PanicException()
```

I'm the 2nd Thread

```
synchronized boolean precombine() {
while (locked) {wait();}
switch (cStatus) {
 case IDLE: cStatus = CStatus.FIRST;
       return true:
 case FIRST: locked = true;
       cStatus = CStatus.SECO
       return false;
 case ROOT: return false;
 default: throw new PanicException
  If 1st thread has promised to return, lock
      node so it won't leave without me
```

Precombining Node

```
synchronized boolean precombine() {
while (locked) {wait();}
switch (cStatus) {
 case IDLE: cStatus = CStatus.FIRST;
        return true;
 case FIRST: locked = true;
        cStatus = CStatus.SECOND;
        return raise;
 case ROOT: return false;
 default: throw new PanicException
                              Prepare to deposit 2<sup>nd</sup>
                                          value
```

Precombining Node

```
synchre
         End of phase 1, don't
while (
            continue up tree
switch
           : cStatus = CStatus.FIRST;
 case ID
       return true;
 case FIR$T: ocked = true;
        cStatus = CStatus.SECOND;
        return false;
 case ROOT: return talse;
 default: throw new PanicException()
```

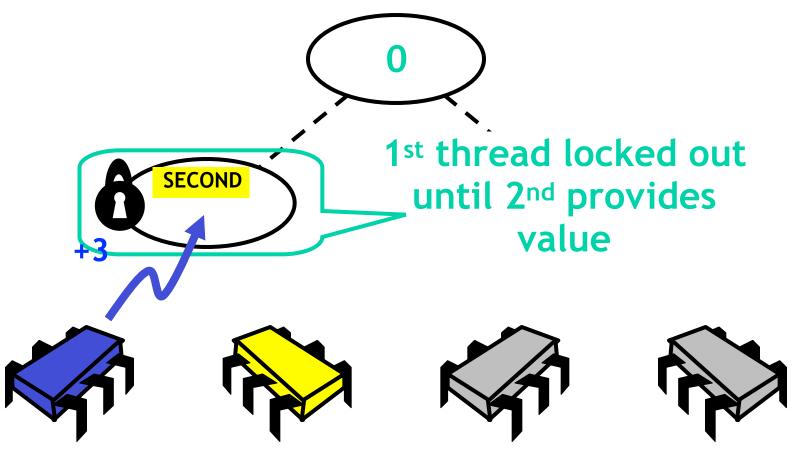
Node is the Root

```
synchronized boolean phase1() {
              If root, phase 1 ends,
while (sStatus
switch (cState don't continue up tree
 case IDLE: c
 case FIRZT: locked = true;
        Status = CStatus.SECOND;
        return false;
 case ROOT: return false;
 default: throw new PanicException()
```

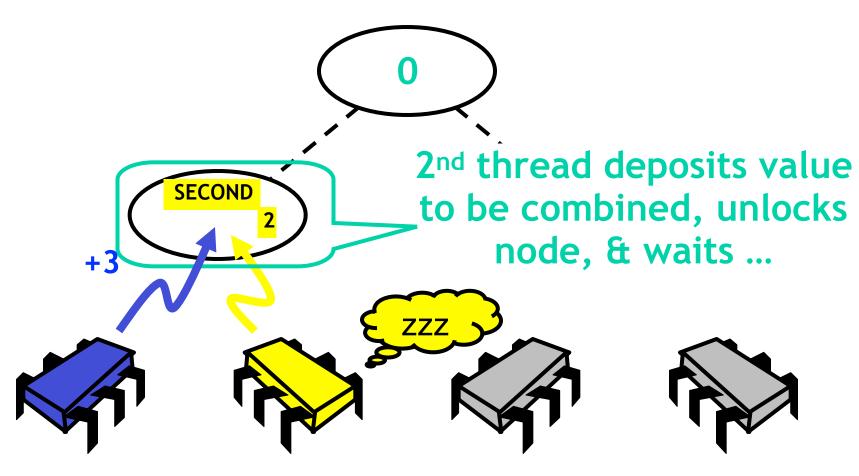
Precombining Node

```
synchronized boolean phase1() {
                  Always check for
while (locked)
switch (cStati unexpected values!
 case IDLE: c
       return true;
 case FIRST: locked = true;
        cStatus = CStatus.SECC
       return false;
 case ROOT: return false;
 default: throw new PanicException()
```

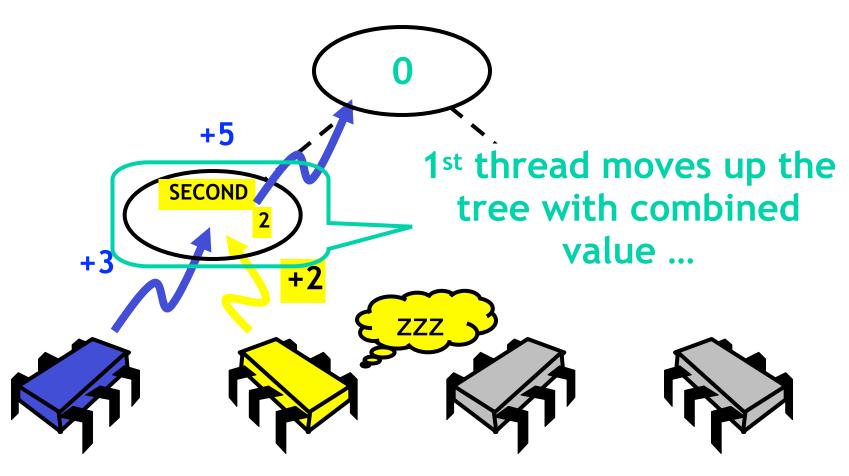
Combining Phase

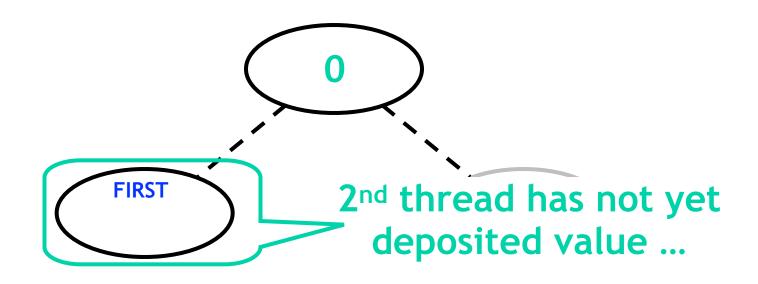


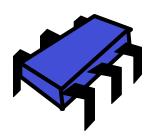
Combining Phase

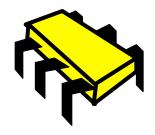


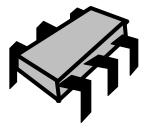
Combining Phase

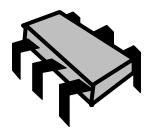


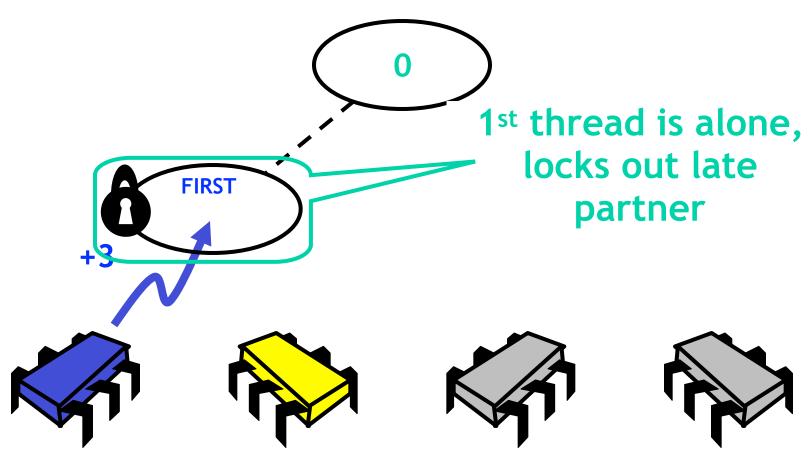




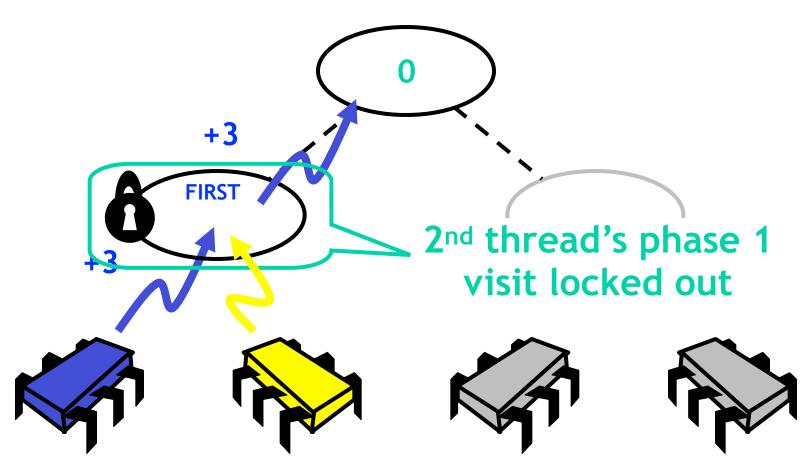








Stop at root +3 **FIRST** +3



```
node = myLeaf;
int combined = 1;
while (node != stop) {
  combined = node.combine(combined);
  stack.push(node);
  node = node.parent;
}
```

```
node = myLeaf;
Int combined = 1,
while (node != stop) {
  combined = node combine(combined);
  stack.push(node);
  node = node.parent;
}
```

Start at leaf

```
node = myLeaf;
int combined = 1;
while (node != stop) {
  combined = node.combine(combined);
  stack.push(node);
  node = node.parent;
}
```

Add 1

```
node = myLeaf;
int combined = 1;
while (node != stop) {
  combined = node.sombine(combined);
  stack.push(node);
  node = node.parent;
}
Revisit nodes visited
in phase 1
```

```
node = myLeaf;
int combined = 1;
while (node != stop) {
  combined = node.combine(combined);
  stack.push(node);
  node = node.parent;
}
```

Accumulate combined values, if any

```
node = myLeaf;
int combined = 1;
while (node != stop) {
  combined = node.combine(combined);
  stack.push(node);
  node = node.parent;
}
```

```
synchronized int combine(int combined) {
 while (locked) wait();
 locked = true;
 firstValue = combined;
 switch (cStatus) {
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue;
  default: ...
```

```
synchronized int combine(int combined) {
 while (locked) wait();
 locked = true
 firstValue = combined
 switch (cStatus) {
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue;
  default: ....
                           Wait until node is unlocked
```

```
synchronized int combine(int combined) {
 while (locked) wait();
 locked = true;
 firstValue = combined
                                 Lock out late attempts
 switch (cStatus) {
                                        to combine
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue;
  default: ....
```

```
synchronized int combine(int combined) {
 while (locked) wait();
 locked = true:
 firstValue = combined;
 switch (cStatus)
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue
  default: ....
                           Remember our contribution
```

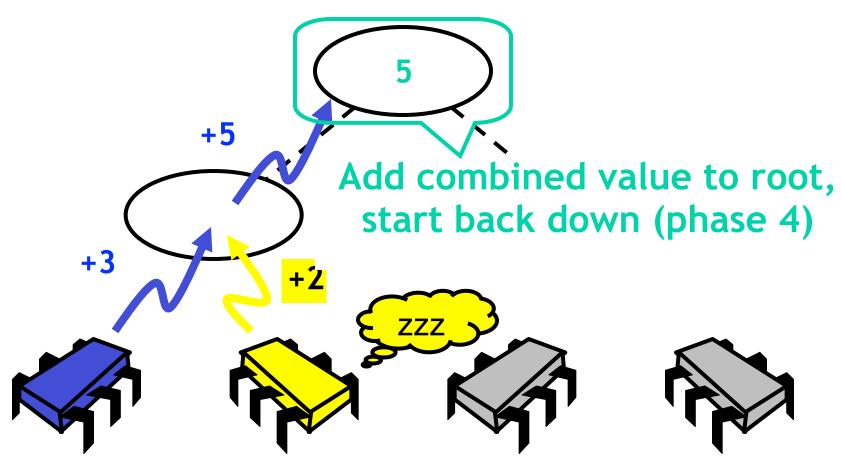
```
synchronized int combine(int combined) {
 while (locked) wait();
                                      Check status
 locked = true;
 firstValue - combi
 switch (cStatus) {
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue;
  default: ....
```

```
synchronized int combine(int combined) {
 while (locked) wait();
 locked = true;
                                  1st thread is alone
 firstValue = combined;
 switch (cStatus)
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue;
  default: ...
```

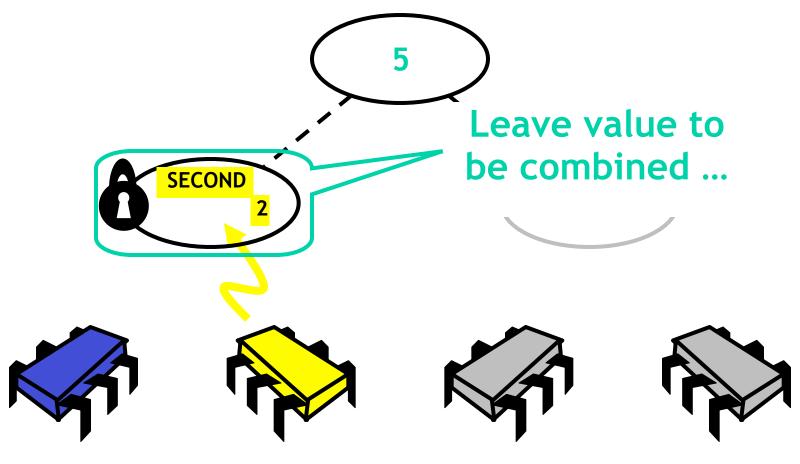
Combining Node

```
synchronized int combine(int combined) {
 while (locked) wait();
 locked = true;
                                      Combine with
 firstValue = combined;
 switch (cStatus) {
                                        2<sup>nd</sup> thread
  case FIRST:
   return firstValue;
  case SECOND:
   return firstValue + secondValue;
  default: ...
```

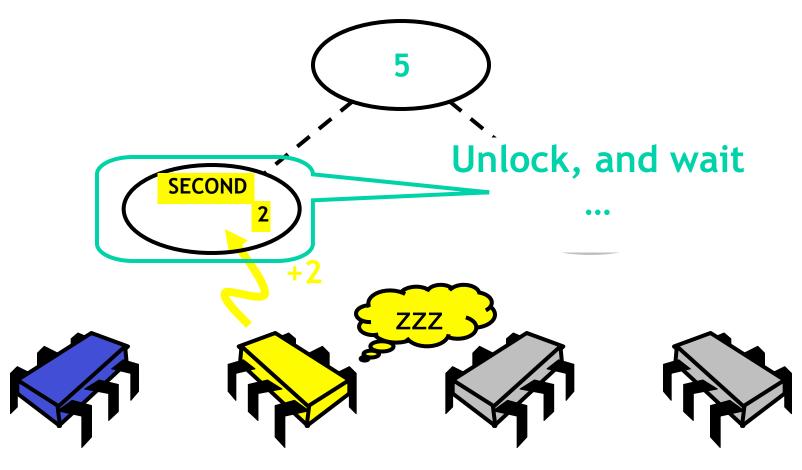
Operation Phase



Operation Phase (reloaded)



Operation Phase (reloaded)



Operation Phase Navigation

prior = stop.op(combined);

Operation Phase Navigation

prior = stop.op(combined); Get result of combining

Operation Phase Node

```
synchronized int op(int combined) {
 switch (cStatus) {
  case ROOT: int oldValue = result;
   result += combined;
   return old Value;
  case SECOND: secondValue = combined;
    locked = false; notifyAll();
   while (cStatus != CStatus.DONE) wait();
    locked = false; notifyAll();
   cStatus = CStatus.IDLE;
   return result;
  default: ....
```

At Root

```
synchronized int op(int combined) {
 switch (cStatus) {
  case ROOT: int oldValue = result;
   result += combined;
   return old Value;
  case SECOND: secondValue = combined;
   locked = false; notifyAll();
   while (cStatus != CStatus.DC
                                    wa t();
   locked = false; notifyAll();
   cStatus = CStatus.IDLE;
                                     Add sum to root,
   return result;
  default: ....
                                     return prior value
```

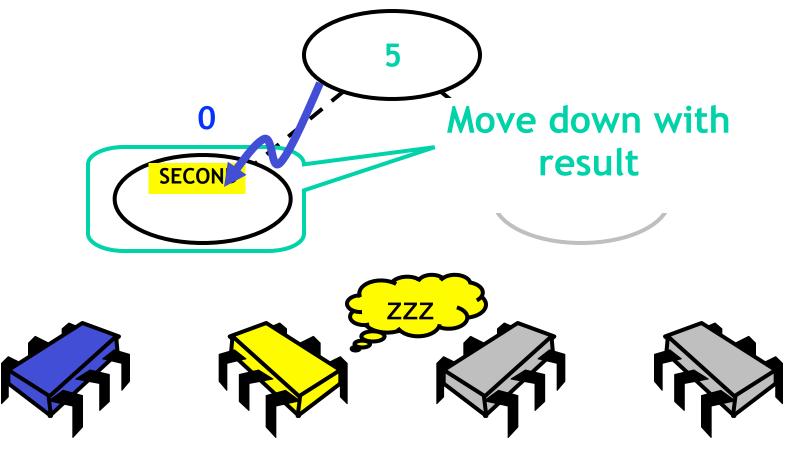
```
synchronized int op(int combined) {
 switch (cStatus) {
  case ROOT: int oldValue = result;
   result += combined;
   return oldValue;
  case SECOND: secondValue = combined;
   locked = false; notifyAll();
   while (cStatus != CStatus.DONE) wait()
   locked = false; notifyAll();
   cStatus = CStatus.IDLE;
                                  Deposit value for later
   return result;
  default: ....
                                         combining ...
```

```
synchronized int op(int combined) {
 switch (cStatus) {
  case ROOT: int oldValue = result;
   result += combined;
   return oldValue;
  case SECOND: second Value - combined:
   locked = false; notifyAll();
   while (cStatus != CStatus.DONE) wait();
   locked = false; notifyAll();
   cStatus = CStatus.IDLE;
                                    Unlock node, notify
   return result;
  default: ...
                                           1st thread
```

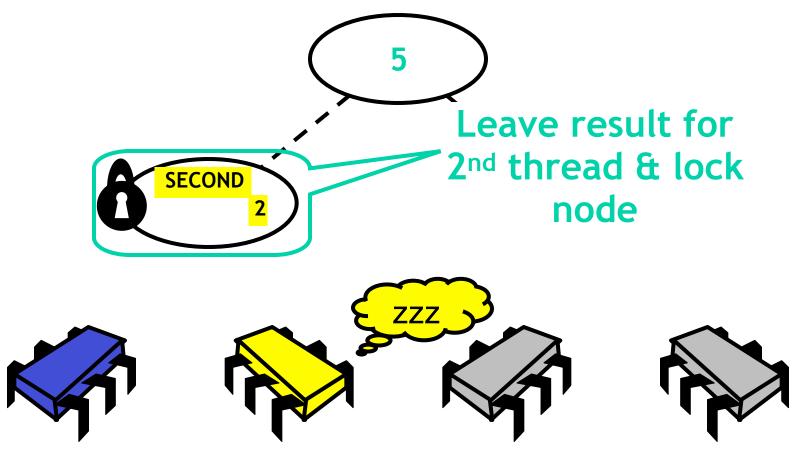
```
synchronized int op(int combined) {
                                        Wait for 1st thread
 switch (cStatus) {
                                         to deliver results
  case ROOT: int oldValue = result;
   result += combined:
   return oldValue;
  case SECOND: secondValue = combin
   locked = false: notifyAll():
   while (cStatus != CStatus.DONE) wait();
   locked = false; notifyAll();
   cStatus = CStatus.IDLE;
   return result;
  default: ....
```

```
synchronized int op(int combined) {
 switch (cStatus) {
                                        Unlock node &
  case ROOT: int oldValue = result;
                                              return
   result += combined;
   return oldValue;
  case SECOND: secondValue = comp
   locked = false; notifyAll();
   while (cStatus != CStatus.DONE
   locked = false; notifyAll();
   cStatus = CStatus.IDLE;
   return result;
  default: ...
```

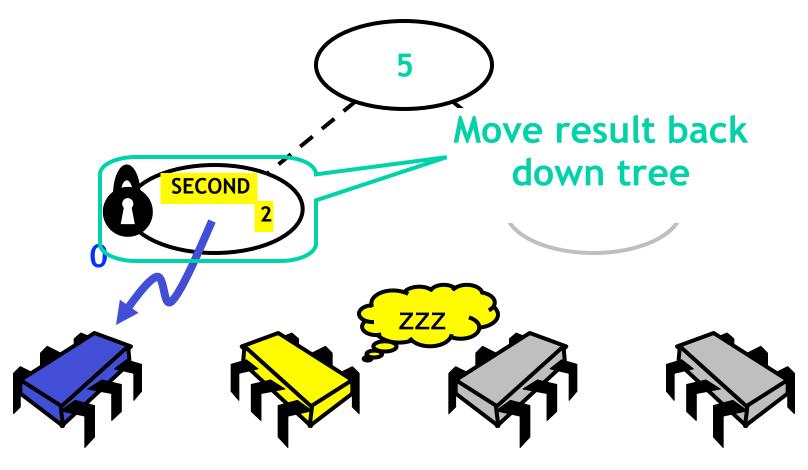
Distribution Phase

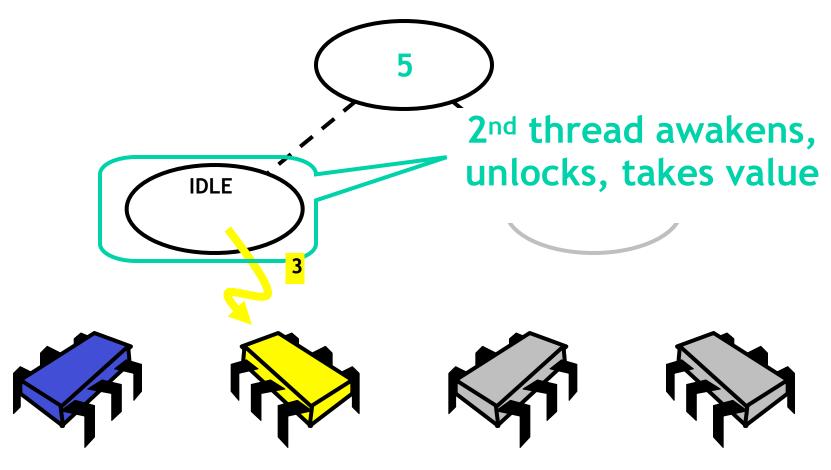


Distribution Phase



Distribution Phase





```
while (!stack.empty()) {
  node = stack.pop();
  node.distribute(prior);
  }
return prior;
```

```
while (!stack.empty()) {
 node = stack.pop();
 node.distribute(pr
return prior;
                          Traverse path in
                           reverse order
```

```
while (!stack.empty()) {
 node - stack pop();
 node.distribute(prior);
return prior;
                          Distribute results to
                          waiting 2<sup>nd</sup> threads
```

```
while (!stack.empty()) {
  node = stack.pop();
  node.distribute(prior);
  }
return prior;
```

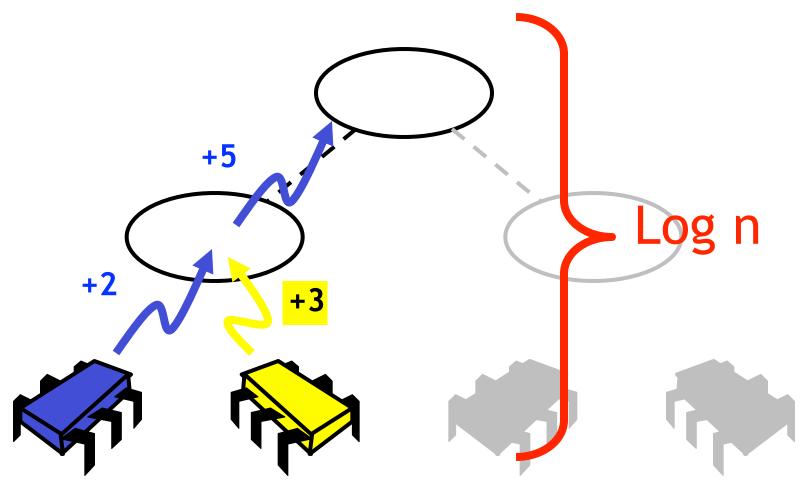
Return result to caller

```
synchronized void distribute(int prior) {
 switch (cStatus) {
   case FIRST:
    cStatus = CStatus.IDLE;
    locked = false; notifyAll();
    return
   case SECOND:
    result = prior + firstValue;
    cStatus = CStatus.DONE; notifyAll();
    return;
   default: ....
```

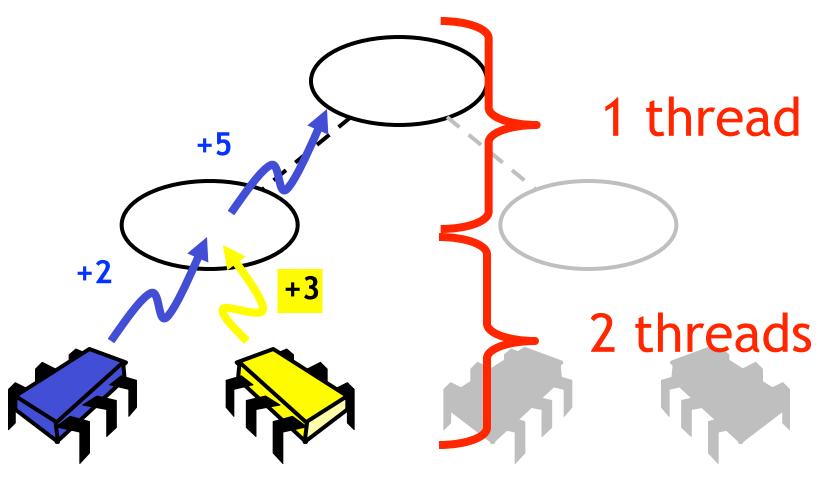
```
synchronized void distribute(int prior) {
 switch (cStatus) {
   case FIRST:
    cStatus = CStatus.IDLE;
    locked = false; notifyAll();
    return;
   case SECOND:
    result = prior + firstValue
    cStatus = CStatus.DONE
    return;
                             No combining, unlock
  default: ....
                                   node & reset
```

```
synchronized void distribute(int prior) {
 switch (cStatus) {
  case FIRS' Notify 2nd thread that
    cStatus = result is available
    locked = 1
    return:
  case SECOND:
    result = prior + firstValue;
    cStatus = CStatus.DONE; notifyAll();
    return;
  default: ...
```

Bad News: High Latency



Good News: Real Parallelism



Throughput Puzzles

- Ideal circumstances
 - All n threads move together, combine
 - n increments in O(log n) time
- Worst circumstances
 - All n threads slightly skewed, locked out
 - n increments in O(n ⋅ log n) time

```
void indexBench(int iters, int work) {
  while (int i < iters) {
    i = r.getAndIncrement();
    Thread.sleep(random() % work);
  }}</pre>
```

```
void indexBench(int iters, in) work) {
  while (int i < iters) {
    i = r.getAndIncrement();
    Thread.sleep(random() % work);
  }}</pre>
```

How many iterations

```
void indexBench(int iters, int work) {
  while (int i < iters) {
    i = r.getAndIncrement(),
    Thread.sleep(random() % work);
  }}</pre>
```

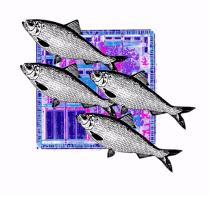
Expected time between incrementing counter

```
void indexBench(int iters, int work) {
while (int i < iters) {
 i = r.getAndIncrement();
 Thread.sleep(random() % work);
                    Pretend to work
           (more work, less concurrency)
```

Performance Benchmarks

- Alewife
 - NUMA architecture
 - Simulated

MIT - ALEWIFE



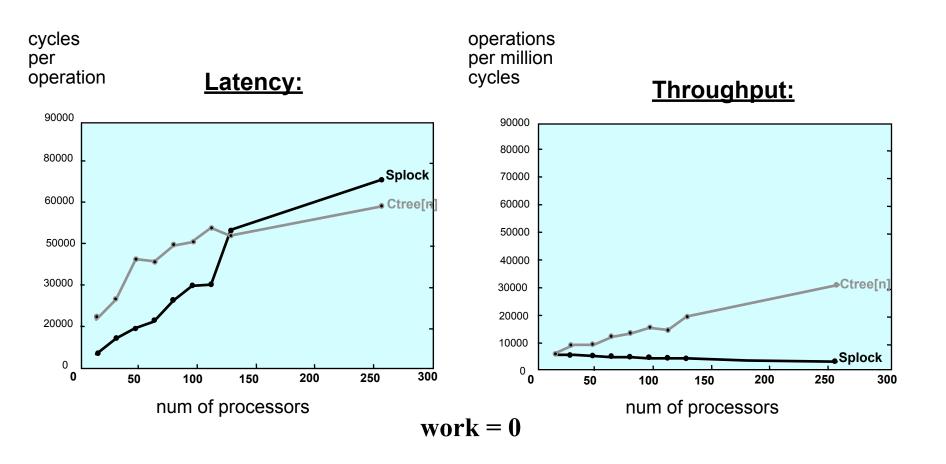
Throughput:

 average number of inc operations in 1 million cycle period.

Latency:

 average number of simulator cycles per inc operation.

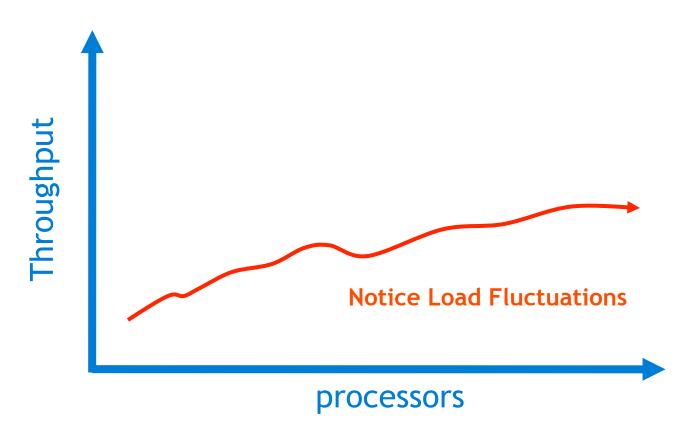
Performance



The Combining Paradigm

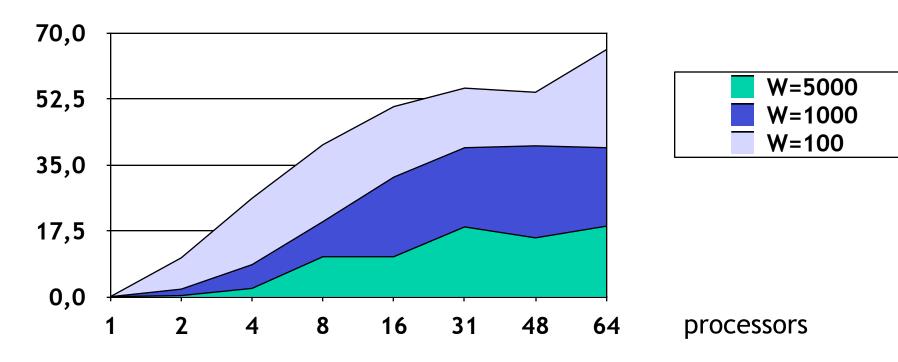
- Implements any RMW operation
- When tree is loaded
 - Takes 2 log n steps
 - for n requests
- Very sensitive to load fluctuations:
 - if the arrival rates drop
 - the combining rates drop
 - overall performance deteriorates!

Combining Load Sensitivity



Combining Rate vs Work

Throughput



Conclusions

- Combining Trees
 - Work well under high contention
 - Sensitive to load fluctuations
 - Can be used for getAndMumble() ops