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Based on slides by Maurice Herlihy

and Nir Shavit

Simple Video Game

- Prepare frame for display
 - By graphics coprocessor
- "soft real-time" application
 - Need at least 35 frames/second
 - OK to mess up rarely

Simple Video Game

```
while (true) {
  frame.prepare();
  frame.display();
}
```

Simple Video Game

```
while (true) {
  frame.prepare();
  frame.display();
  }
```

- What about overlapping work?
 - 1st thread displays frame
 - 2nd prepares next frame

Two-Phase Rendering

```
while (true) {
 if (phase) {
  frame[0].display();
 } else {
  frame[1].display();
 phase = !phase;
```

```
while (true) {
 if (phase) {
  frame[1].prepare();
 } else {
  frame[0].prepare();
 phase = !phase;
```

Two-Phase Rendering

```
while (true) {
                            while (true) {
 if (phase) {
                              if (phase) {
  frame[0].display();
                               frame[1].prepare();
 } else {
                              } else {
  frame[1].display();
                               frame[0].prepare();
 phase = !phase;
                                    =!phase;
```

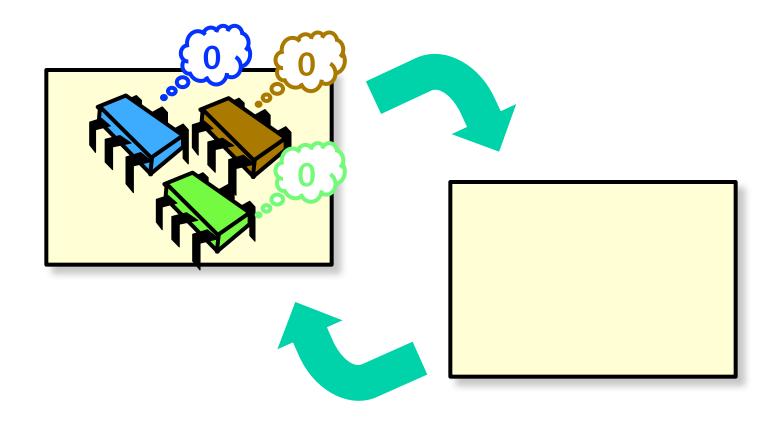
Even phases

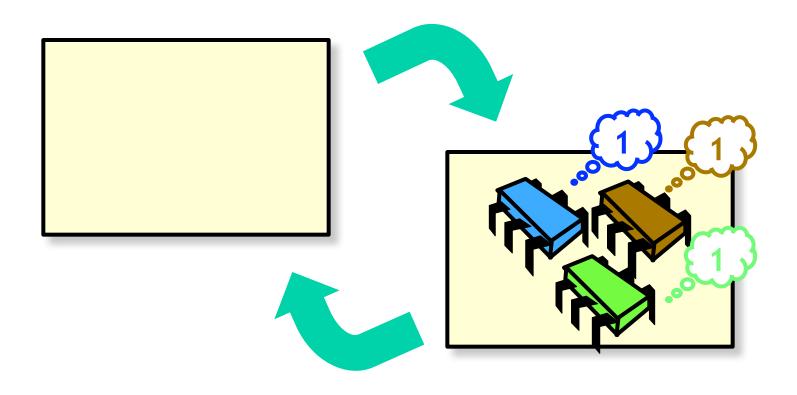
Two-Phase Rendering

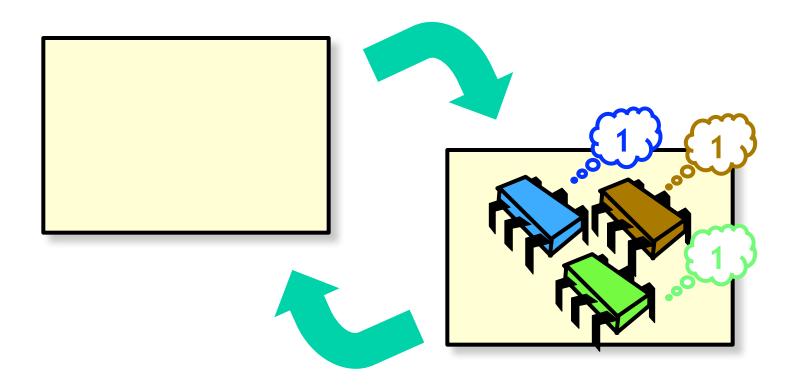
```
while (true) {
                            while (true) {
 if (phase) {
                             if (phase) {
  frame[0].display();
                               frame[1].prepare();
 } else {
                             } else {
  frame[1].display();
                               frame[0].prepare();
 phase = !phase;
                                nase =
                    odd phases
```

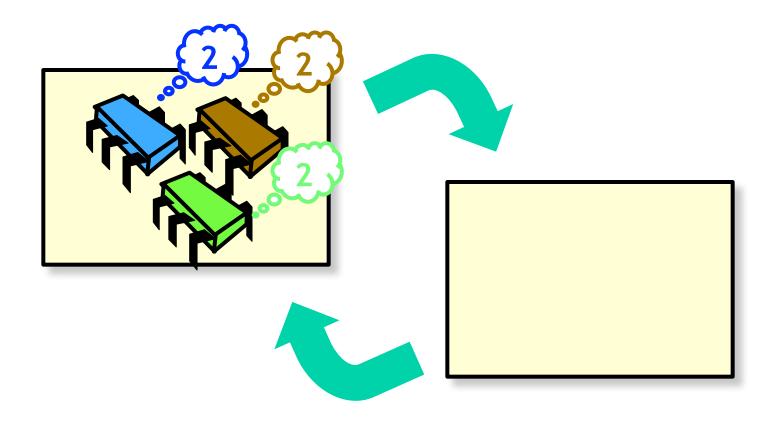
Synchronization Problems

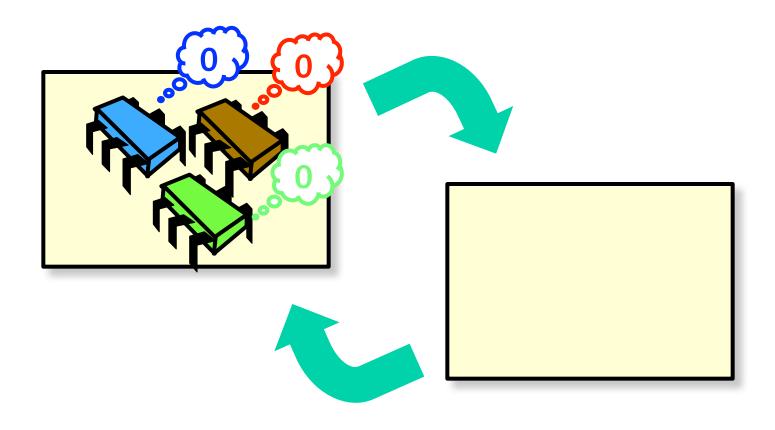
- How do threads stay in phase?
- Too early?
 - "we render no frame before its time"
- Too late?
 - Recycle memory before frame is displayed

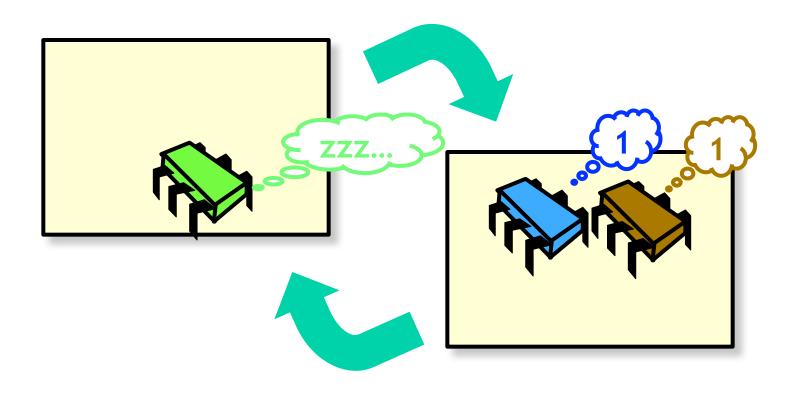


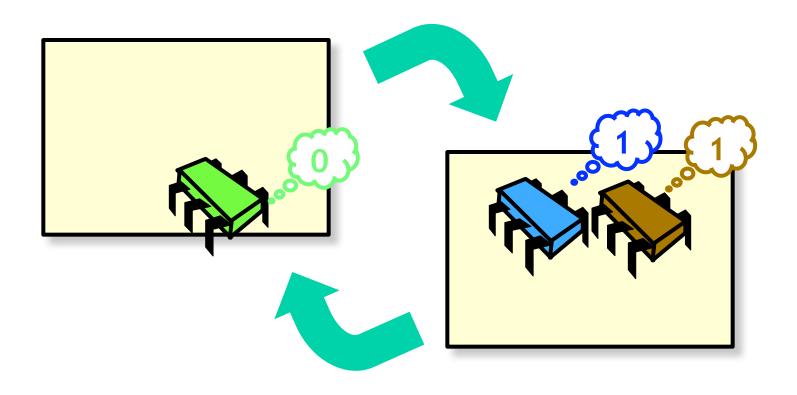


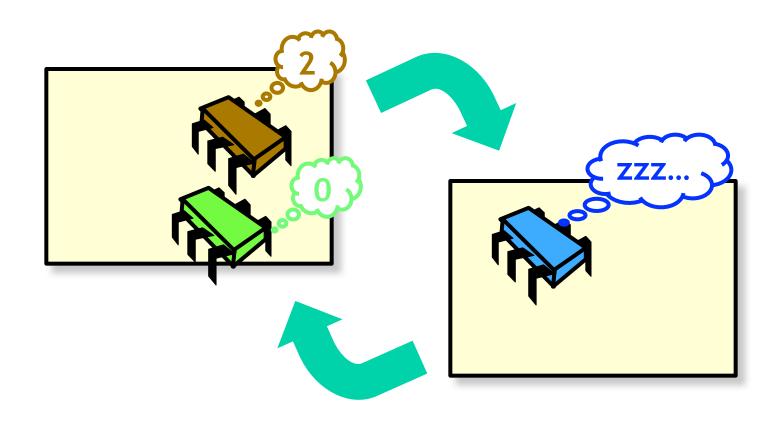


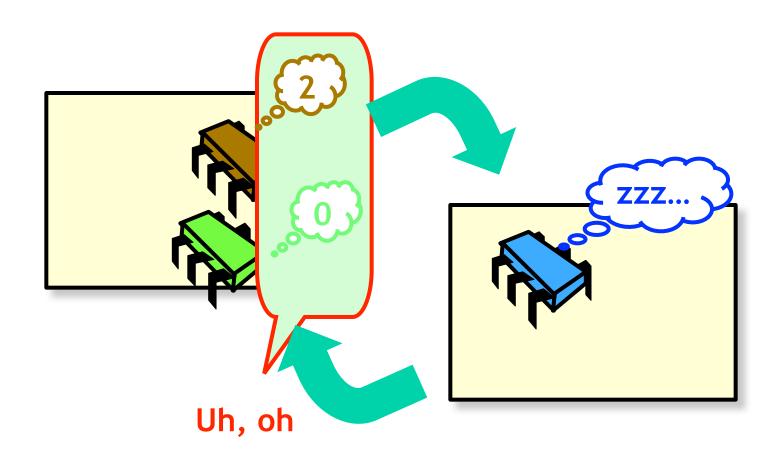


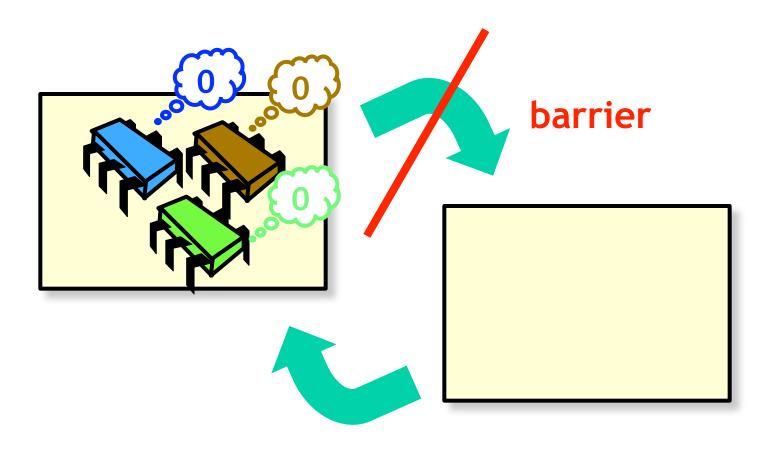


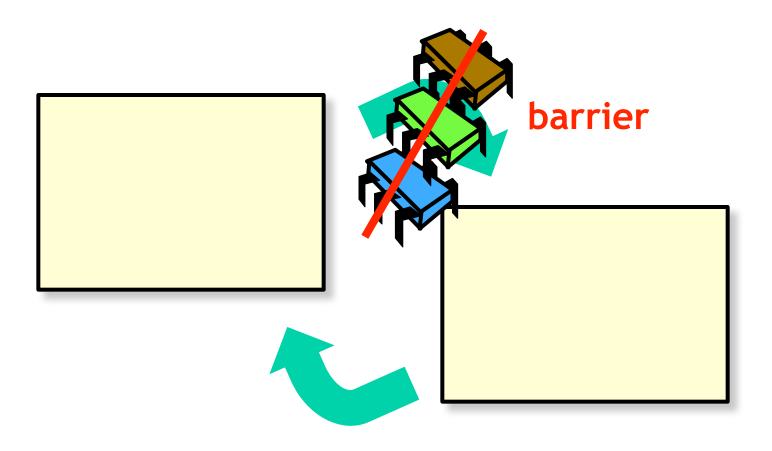


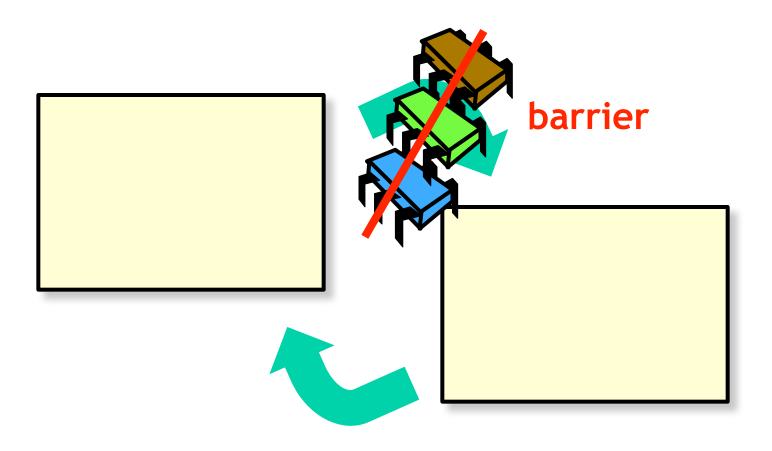


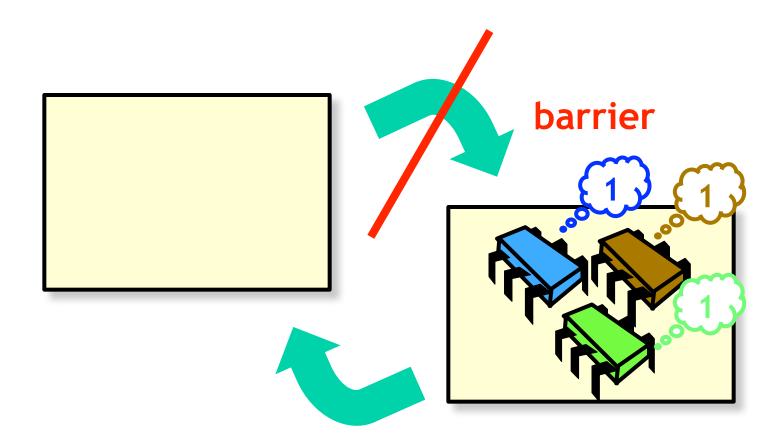


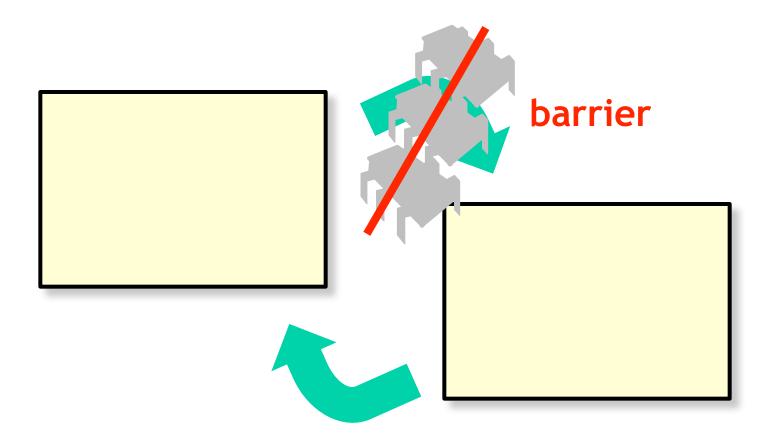


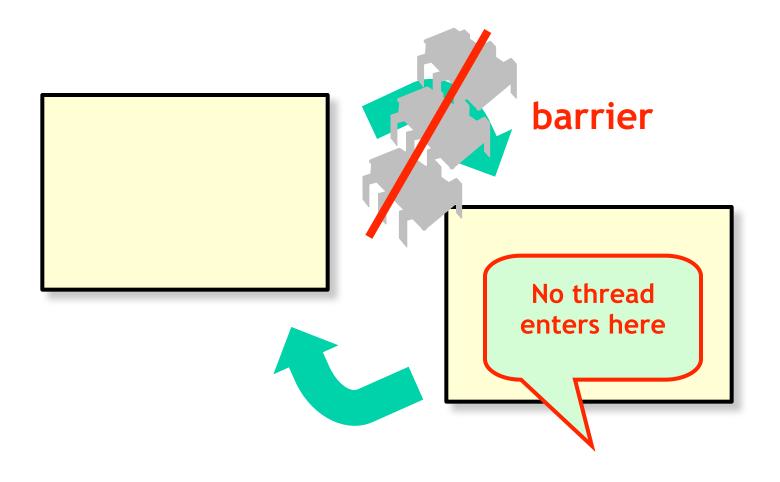


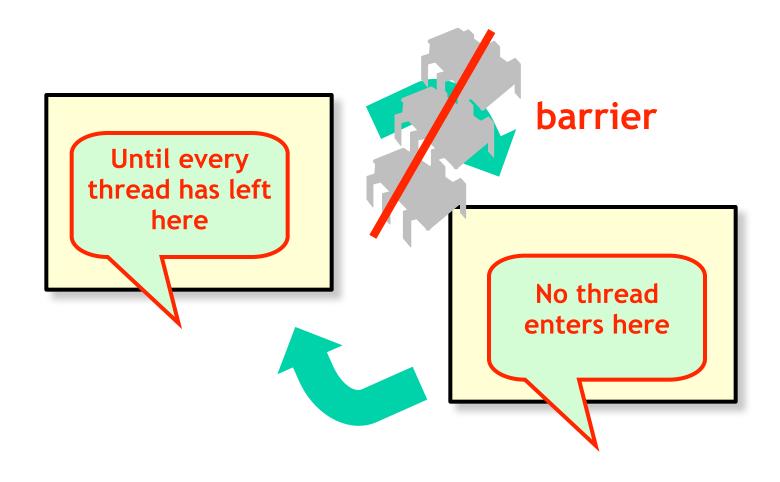












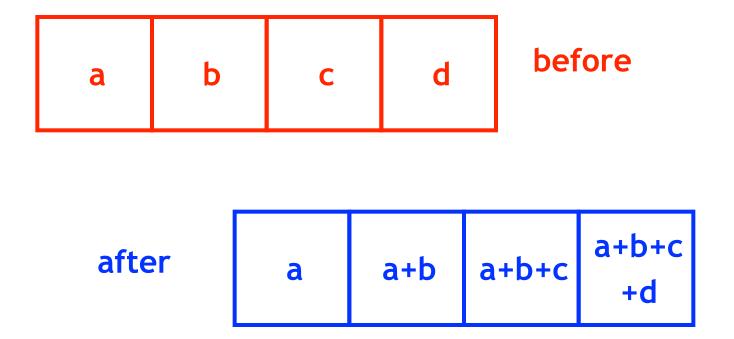
Why Do We Care?

- Mostly of interest to
 - Scientific & numeric computation
- Elsewhere
 - Garbage collection
 - Less common in systems programming
 - Still important topic

Duality

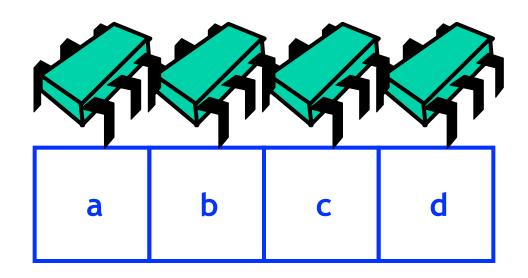
- Dual to mutual exclusion
 - Include others, not exclude them
- Same implementation issues
 - Interaction with caches ...
 - Invalidation?
 - Local spinning?

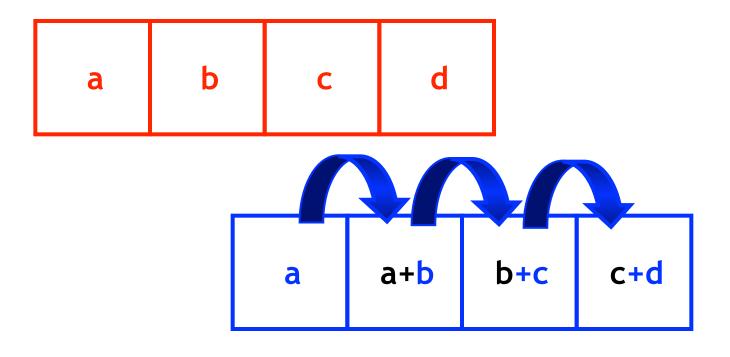
Example: Parallel Prefix

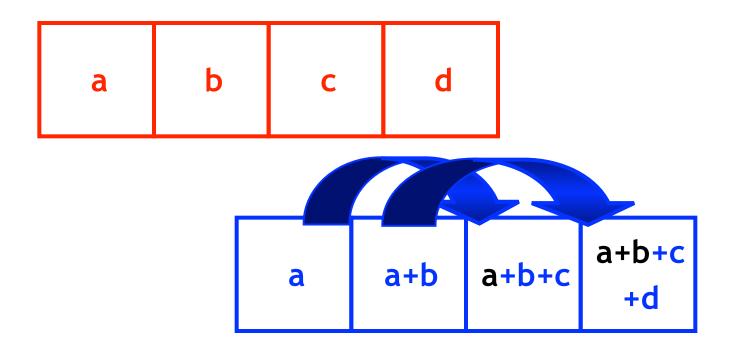


Parallel Prefix

One thread Per entry

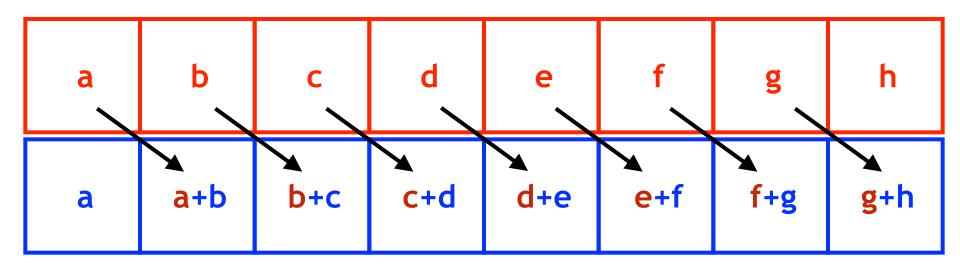




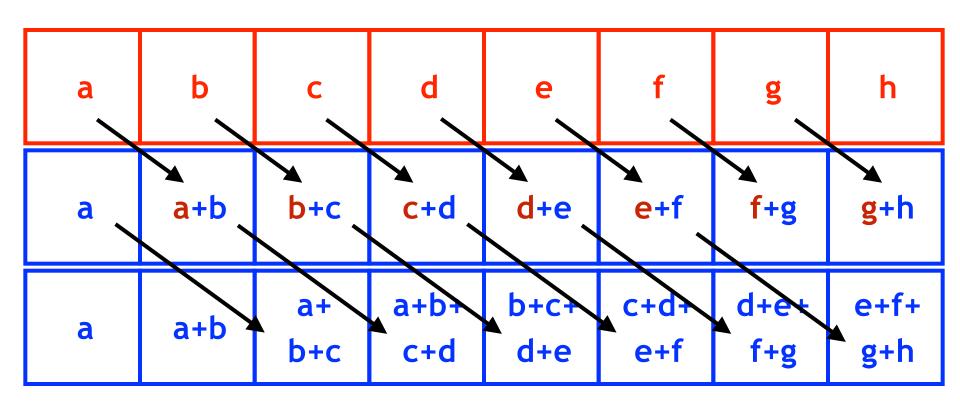


a b c d e f g h

a	b	С	d	е	f	g	h
a	a+b	b+c	c+d	d+e	e+f	f+g	g+h

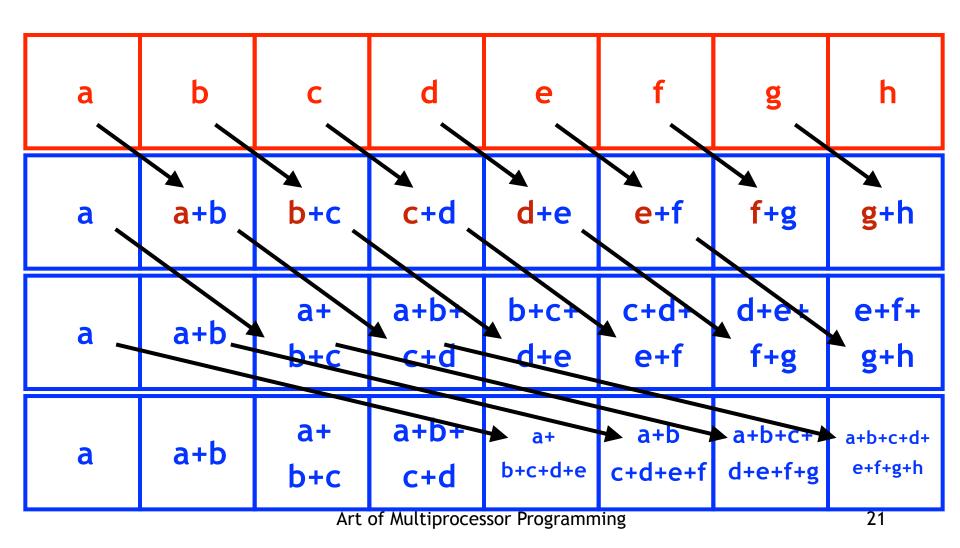


a	b	С	d	е	f	g	h
a	a+b	b+c	c+d	d+e	e+f	f+g	g+h
a	a+b	a+ b+c	a+b+ c+d	b+c+ d+e	c+d+ e+f	d+e+ f+g	e+f+ g+h



a	b	С	d	е	f	g	h
a	a+b	b+c	c+d	d+e	e+f	f+g	g+h
a	a+b	a+ b+c	a+b+ c+d	b+c+ d+e	c+d+ e+f	d+e+ f+g	e+f+ g+h
a	a+b	a+ b+c	a+b+ c+d	a+ b+c+d+e ssor Programm	a+b c+d+e+f	a+b+c+ d+e+f+g	a+b+c+d+ e+f+g+h 21

Parallel Prefix: Phase 2



Parallel Prefix

- N threads can compute
 - Parallel prefix
 - Of N entries
 - In log₂ N rounds
- What if system is asynchronous?
 - Why we need barriers

```
class Prefix extends Thread {
private int[] a;
private int i;
private Barrier b;
public Prefix(int[] a,
          Barrier b, int i) {
 this.a = a;
 this.b = b;
 this.i = i;
```

```
class Prefix extends Thread {
private int[] a;
private int i;
private Barrier b;
public Prefix(int[] a,
          Barrier b, int i) {
 this.a = a;
                             Array of input values
 this.b = b;
this.i = i;
```

```
class Prefix extends Thread {
private int[] a;
private int i;
private Barrier
public Prefix(int[] a,
          Barrier b, int i)
 this.a = a;
                                   Thread index
 this.b = b;
this.i = i;
```

```
class Prefix extends Thread {
private int[] a;
private int i;
private Barrier b;
public Prefix(int[] a,
          Barrier b, int
 this.a = a;
                            Shared barrier
 this.b = b;
this.i = i;
```

```
class Prefix extends Thread {
private int[] a;
private int i;
                                     Initialize fields
private Barrier b;
public Prefix(int[] a,
          Barrier b, int j
 this.a = a;
 this.b = b;
 this.i = i;
```

```
public void run() {
int d = 1, sum = 0;
while (d < N) {
 if (i >= d)
  sum = a[i-d];
 if (i >= d)
 a[i] += sum;
 d = d * 2;
```

```
public void run() {
int d = 1, sum = 0;
while (d < N) {
 if (i >= d)
  sum = a[i-d];
 b.await();
 if (i >= d)
  a[i] += sum;
 d = d * 2;
}}}
```

```
public void run() {
int d = 1, sum = 0;
while (d < N) {
 if (i >= d)
  sum = a[i-d];
                     Make sure everyone reads before
 b.await();
                                 anyone writes
 if (i >= d)
 a[i] += sum;
 d = d * 2;
}}}
```

```
public void run() {
int d = 1, sum = 0;
while (d < N) {
 if (i >= d)
  sum = a[i-d];
                     Make sure everyone reads before
 b.await();
                                  anyone writes
 if (i >= d)
 a[i] += sum;
 b.await();
 d = d * 2;
}}}
```

```
public void run() {
int d = 1, sum = 0;
while (d < N) {
 if (i >= d)
  sum = a[i-d];
                    Make sure everyone reads before
 b.await();
                               anyone writes
 if (i >= d)
 a[i] += sum;
                        Make sure everyone writes
 b.await();
                            before anyone reads
 d = d * 2;
}}}
```

Barrier Implementations

- Cache coherence
 - Spin on locally-cached locations?
 - Spin on statically-defined locations?
- Latency
 - How many steps?
- Symmetry
 - Do all threads do the same thing?

```
public class Barrier {
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecrement()==1) {
 count.set(size);
 } else {
 while (count.get() != 0);
}}}}
```

```
public class Barrier {
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
                                  Number threads not
 size = n;
                                       yet arrived
public void await() {
 if (count.getAndDecrement()==1) {
 count.set(size);
} else {
 while (count.get() != 0);
                  Art of Multiprocessor Programming
```

```
public class Barrier {
AtomicInteger count;
                                      Number threads
int size;
                                         participating
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecrement()==1) {
  count.set(size);
 } else {
 while (count.get() != 0);
                  Art of Multiprocessor Programming
```

```
public class Barrier {
                                            Initialization
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecrement()==1) {
  count.set(size);
 } else {
 while (count.get() != 0);
                   Art of Multiprocessor Programming
```

```
Principal method
public class Barrier {
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecrement()==1) {
  count.set(size);
 } else {
 while (count.get() != 0);
                  Art of Multiprocessor Programming
```

```
public class Barrier {
                                If I'm last, reset fields
AtomicInteger count;
                                      for next time
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await()
 if (count.getAndDecrement()==1) {
  count.set(size);
 } else {
  while (count.get() != 0);
                  Art of Multiprocessor Programming
```

```
Otherwise, wait for
public class Barrier {
                                     everyone else
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecreptent()==1) {
 count.set(size);
 } else {
  while (count.get() != 0);
```

```
public class Barrier {
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
       What's wrong with this protocol?
public void await() {
 if (count.getAndDecrement()==1) {
 count.set(size);
 } else {
 while (count.get() != 0);
```

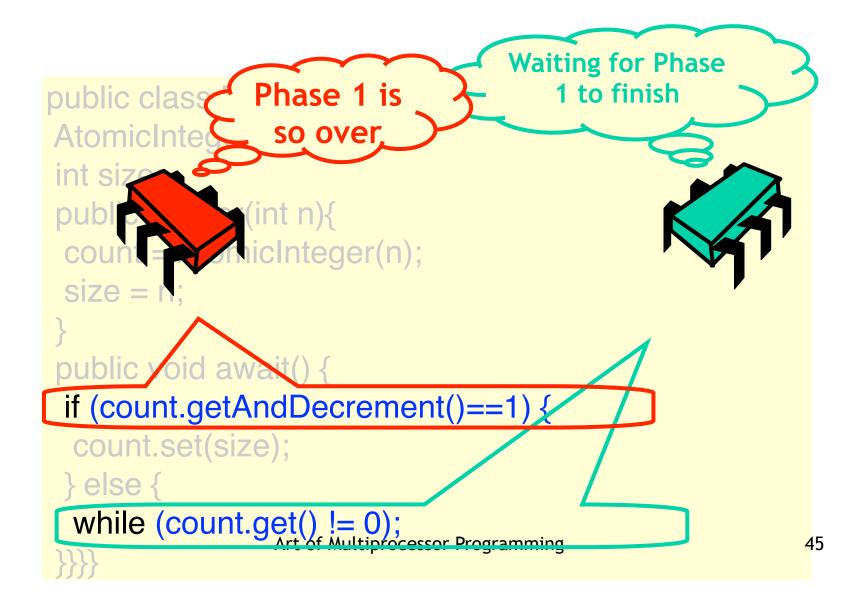
```
Barrier b = new Barrier(n);
while ( mumble() ) {
  work();
  b.await()
}
```

```
Barrier b = new Barrier(n);
while ( mumble() ) {
    work();
    Do work
    synchronize
```

```
Barrier b = new Barrier(n);
while ( mumble() ) {
    work();
    Do work
    synchronize
```

```
public class Barrier {
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecrement()==1) {
 count.set(size);
 } else {
 while (count.get() != 0);
```

```
Waiting for Phase
public class Barrier {
                                     1 to finish
AtomicInteger count;
int size;
public Barrier(int n){
 count = AtomicInteger(n);
 size = n;
public void await() {
 if (count.getAndDecrement()==1)
 count.set(size);
 } else {
 while (count.get() != 0);
```



```
ZZZZZ.
                 Prepare for
public class
                   phase 2
AtomicInte
              er(int n){
            tomicInteger(n);
 size =
        void await() {
    (count.getAndDecrement()==
 count.set(size);
   else
 while (count.get() != 0);
Art of Multiprocessor Programming
                                                                46
```

Uh-Oh

```
Waiting for Phase
                              2 to finish
public class Barrier
                                       ing for Phase
                                      1 to finish
            rier(int n){
            tomicInteger(n);
 count =
 size =
public void await() {
 if (count.getAndDecrement()==1
 count.set(size);
  while (count.get(), != 0);
                  Art or mattiprocessor i rogramming
```

Basic Problem

- One thread "wraps around" to start phase 2
- While another thread is still waiting for phase 1
- One solution:
 - Always use two barriers

```
public class Barrier {
AtomicInteger count;
int size;
volatile boolean sense = false;
threadSense = new ThreadLocal<br/>boolean>...
public void await {
 boolean mySense = threadSense.get();
 if (count.getAndDecrement()==1) {
 count.set(size); sense = mySense
 } else {
 while (sense != mySense) {}
threadSense.set(!mySense)}}}
```

```
Completed odd or
public class Barrier {
                                              even-numbered
AtomicInteger count;
                                                    phase?
int size:
volatile boolean sense = false;
public void await {
 boolean mySense = threadSense.get();
 if (count.getAndDecrement()==1) {
 count.set(size); sense = mySense
 } else {
 while (sense != mySense) {}
threadSense.set(!mySense)}}
```

```
Store sense for next
public class Barrier {
                                               phase
AtomicInteger count;
int size;
volatile boolean sense = false;
threadSense = new ThreadLocal<boolean>...
public void await {
 boolean mySense = threadSense.get();
 if (count.getAndDecrement()==1) {
 count.set(size); sense = mySense
 } else {
 while (sense != mySense) {}
threadSense.set(!mySense)}}
```

```
public class Barrier {
AtomicInteger count;
                                    Get new sense
int size;
                             determined by last phase
volatile boolean sense = false
threadSense = new ThreadLo
public void await {
 boolean mySense = threadSense.get();
 if (count.getAndDecrement()==1) {
 count.set(size); sense = mySense
 } else {
 while (sense != mySense) {}
threadSense.set(!mySense)}}
```

```
public class Barrier {
                              If I'm last, reverse sense
AtomicInteger count;
                                     for next time
int size;
volatile boolean sense = false;
threadSense = new ThreadLocal<br/>boolean>...
public void await {
 boolean mySense = threadSense.get();
 if (count.getAndDecrement()==1) {
 count.set(size); sense = mySense
 } else {
 while (sense != mySense) {}
threadSense.set(!mySense)}}
```

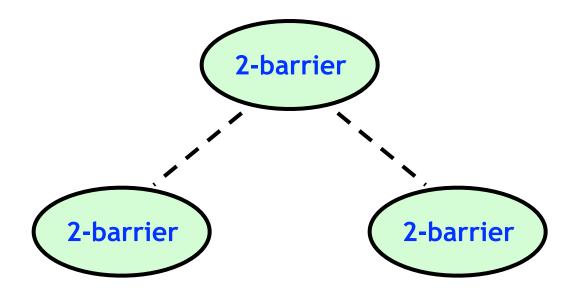
Sense-Reversing Barriers

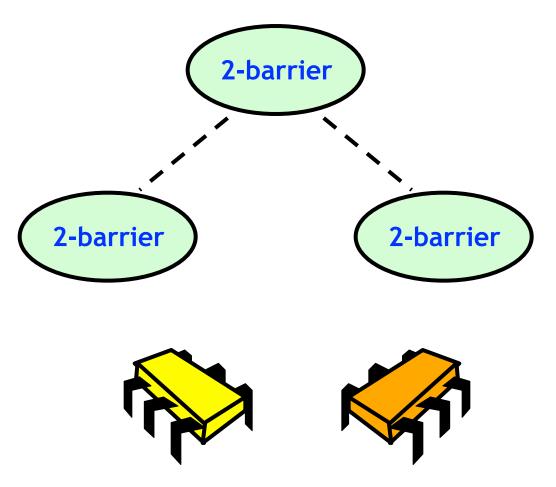
```
public class Barrier {
                              Otherwise, wait for sense
AtomicInteger count;
                                          to flip
int size;
volatile boolean sense = false;
threadSense = new ThreadLocal<boolean>...
public void await {
 boolean mySense = threadSense get()
 if (count.getAndDecrement()==/) {
 count.set(size); sense = my$ense
  else {
 while (sense != mySense) {}
threadSense.set(!mySense)}}
```

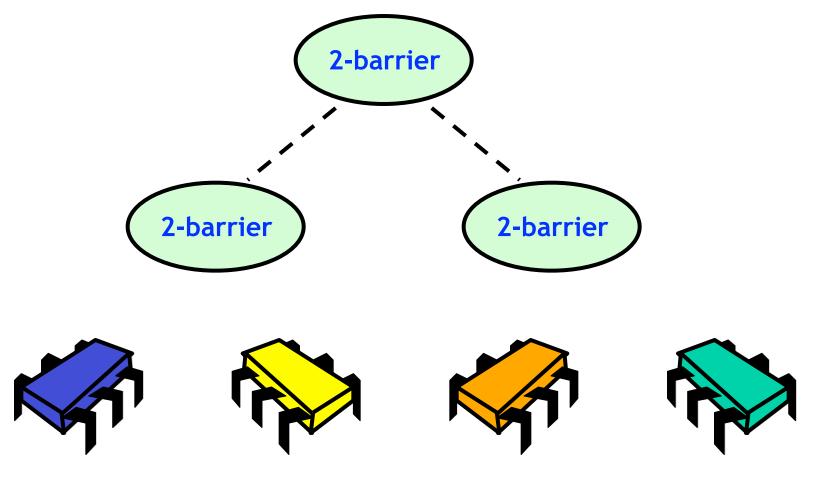
Sense-Reversing Barriers

```
public class Barrier {
                               Prepare sense for next
AtomicInteger count;
                                          phase
int size;
volatile boolean sense = false;
threadSense = new ThreadLocal<boole_n>
public void await {
 boolean mySense = threadSense.get(
 if (count.getAndDecrement()==1
 count.set(size); sense = mySerse
 } else {
 while (sense != mySense)
threadSense.set(!mySense)}}}
```

- Reduce contention
- Split large barrier into tree of small barriers
- Requests go up the tree and notifications down
- Adds latency







```
public class Node{
AtomicInteger count; int size;
Node parent; volatile boolean sense;
public void await() {...
 if (count.getAndDecrement()==1) {
  if (parent != null)
  parent.await()
 count.set(size);
 sense = mySense
 } else {
 while (sense != mySense) {}
}...}}
```

```
Parent barrier in tree
public class Node{
AtomicInteger count: intalze
Node parent; Volatile boolean sense;
public void await() {...
 if (count.getAndDecrement()==1) {
  if (parent != null)
  parent.await()
  count.set(size);
  sense = mySense
 } else {
  while (sense != mySense) {}
}...}}
                    Art of Multiprocessor Programming
```

```
Am I last?
public class Node{
AtomicInteger count; int size,
Node parent; Volatile boolean sep
public void await() {...
 if (count.getAndDecrement()==1) {
  if (parent != null)
  parent.await()
  count.set(size);
  sense = mySense
 } else {
  while (sense != mySense) {}
}...}}
                    Art of Multiprocessor Programming
```

```
public class Node{
                                         Proceed to parent
AtomicInteger count; int size;
                                                barrier
Node parent; Volatile boolean sen
public void await() {...
 if (count_getAndDecre
 if (parent != null)
  parent.await();
 count.set(size);
 sense = mySense
 } else {
 while (sense != mySense) {}
}...}}
```

```
public class Node{
                                      Prepare for next
AtomicInteger count; int size;
                                             phase
Node parent; Volatile boolean sense
public void await() {...
 if (count.getAndDecry
 if (parent != null)
  parent.await()
 count.set(size);
 sense = mySense
 } else {
 while (sense != mySense) {}
}...}}
```

```
public class Node{
                                Notify others at this
AtomicInteger count; int size;
                                          node
Node parent; Volatile boolean
public void await() {..
 if (count.getAndDegreenent()==1) {
 if (parent != null)
  parent.await()
 count.set(size
 sense = mySense
 } else {
 while (sense != mySense) {}
}...}}
```

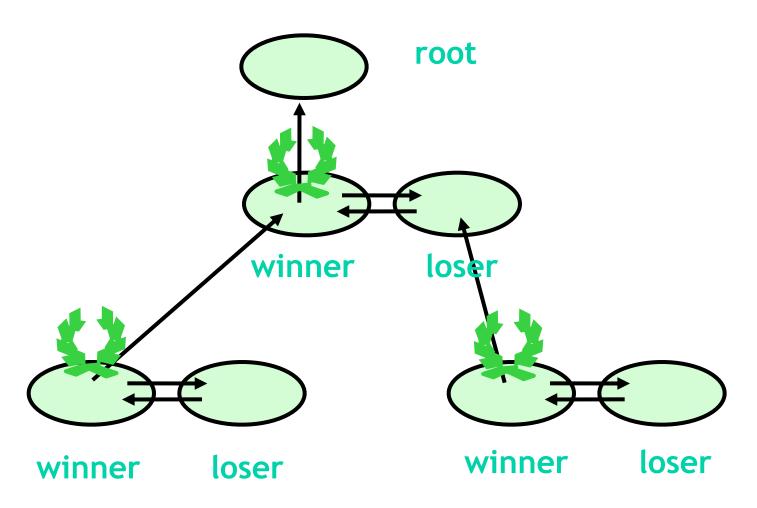
```
public class Node{
                                      I'm not last, so wait
AtomicInteger count; int size;
                                        for notification
Node parent; Volatile boolean sense
public void await() {...
 if (count.getAndDecrement
 if (parent != null) {
  parent.await()
 count.set(size);
 sense = mySense
 } else {
 while (sense != mySense) {}
```

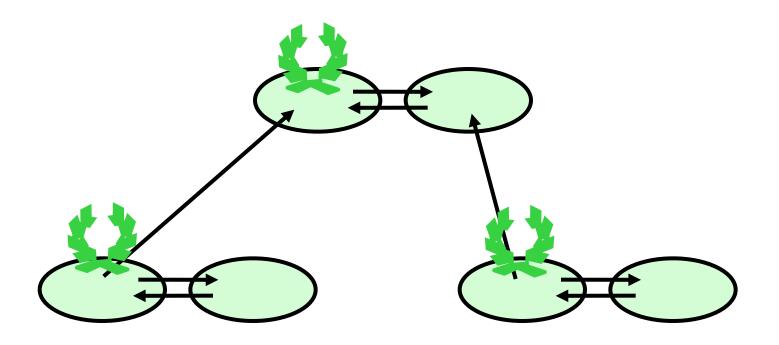
- No sequential bottleneck
 - Parallel getAndDecrement() calls
- Low memory contention
 - Same reason
- Cache behavior
 - Local spinning on bus-based architecture
 - Not so good for NUMA

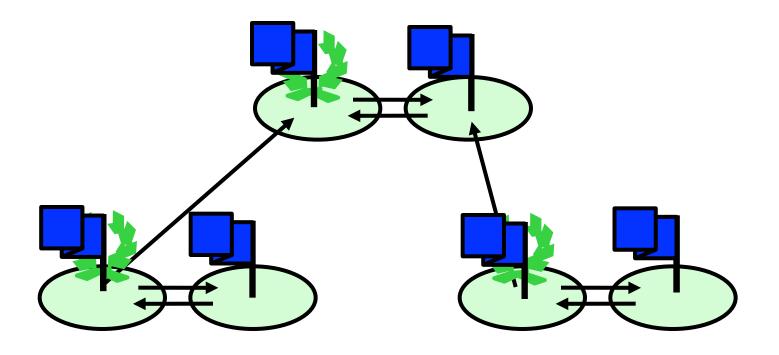
Remarks

- Everyone spins on sense field
 - Local spinning on bus-based (good)
 - Network hot-spot on distributed architecture (bad)
- Not really scalable

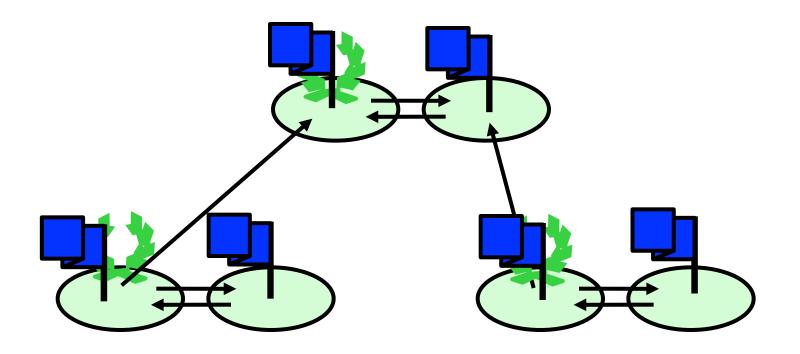
- If tree nodes have fan-in 2
 - Don't need to call getAndDecrement()
 - Winner chosen statically
- At level i
 - If i-th bit of id is 0, move up
 - Otherwise keep back

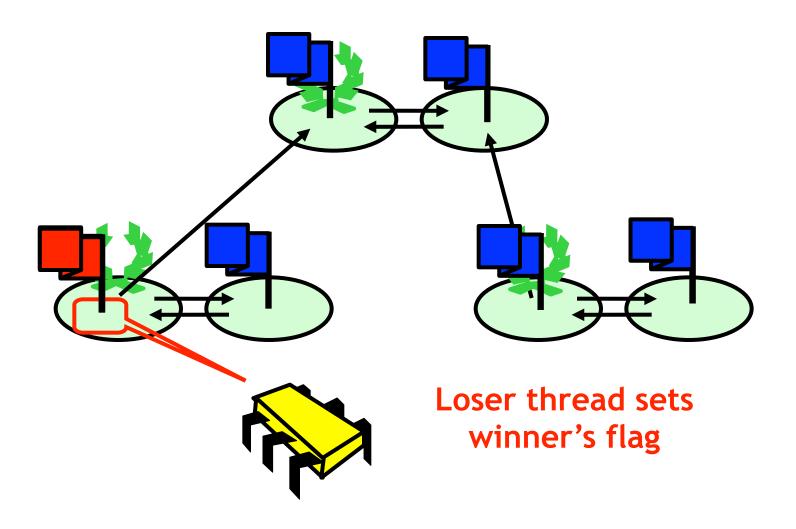


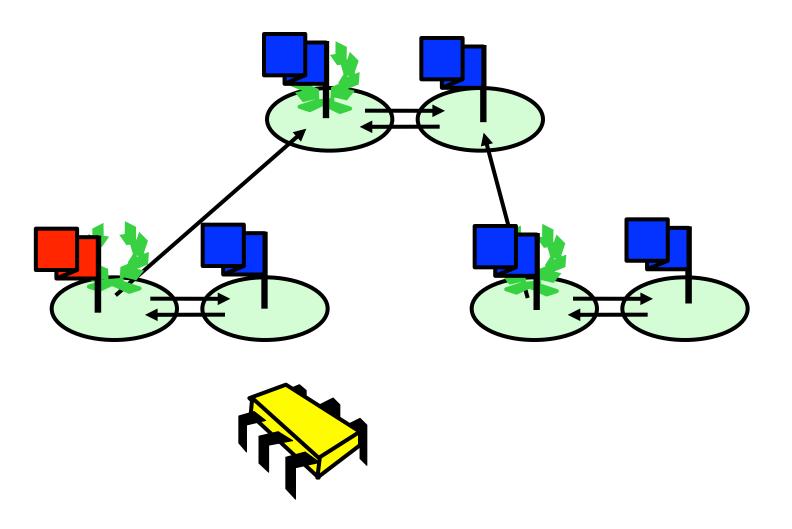


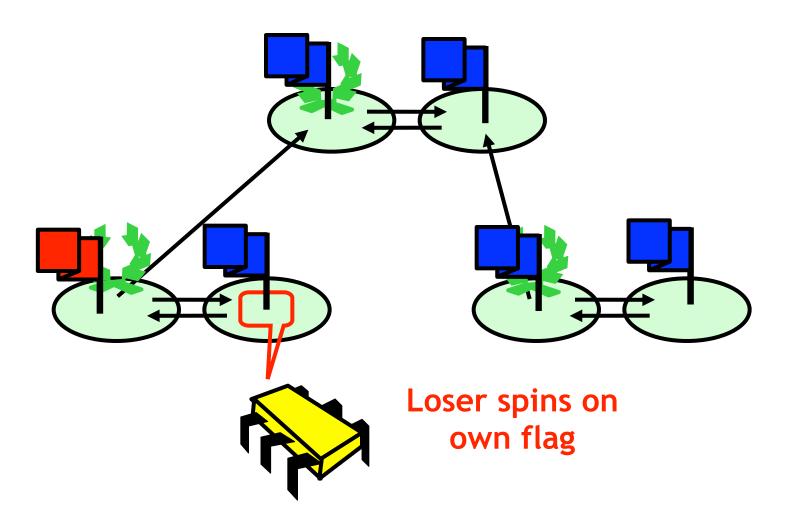


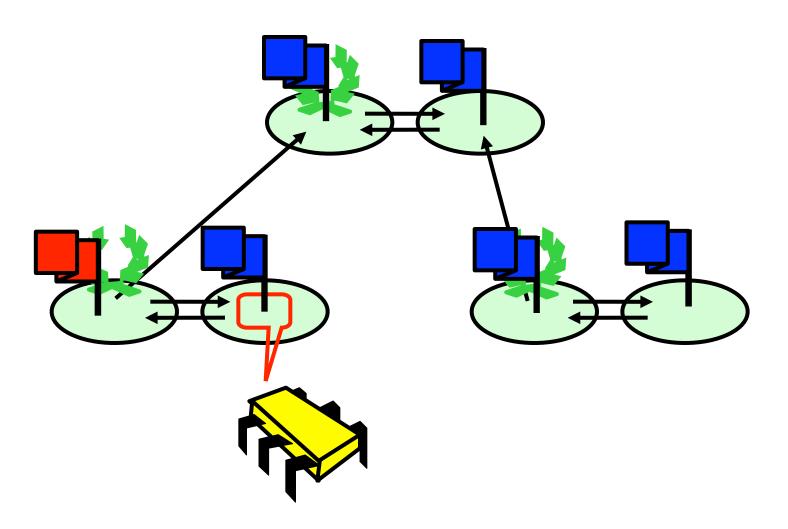
All flags blue

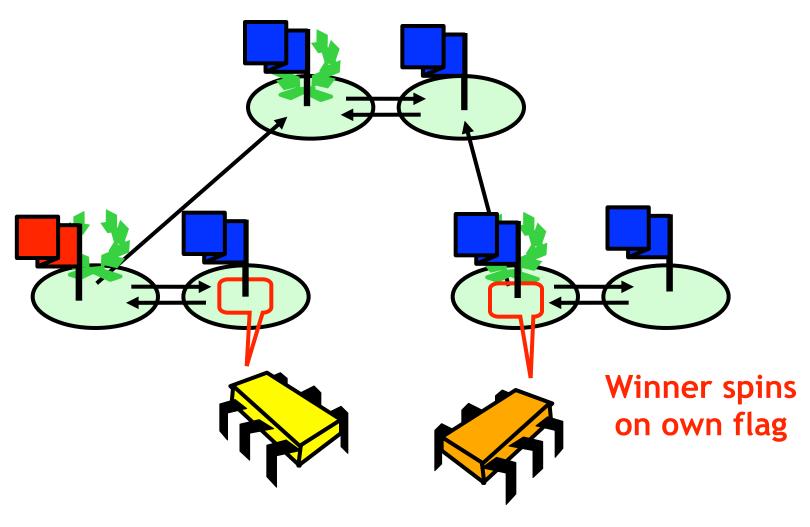


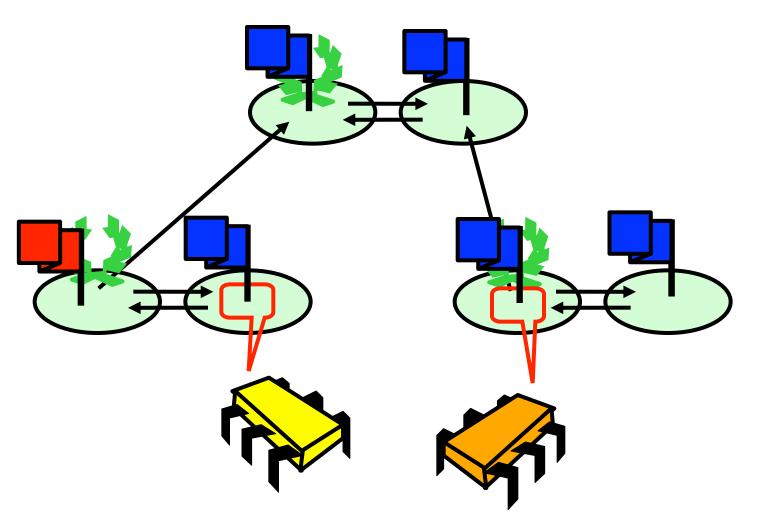


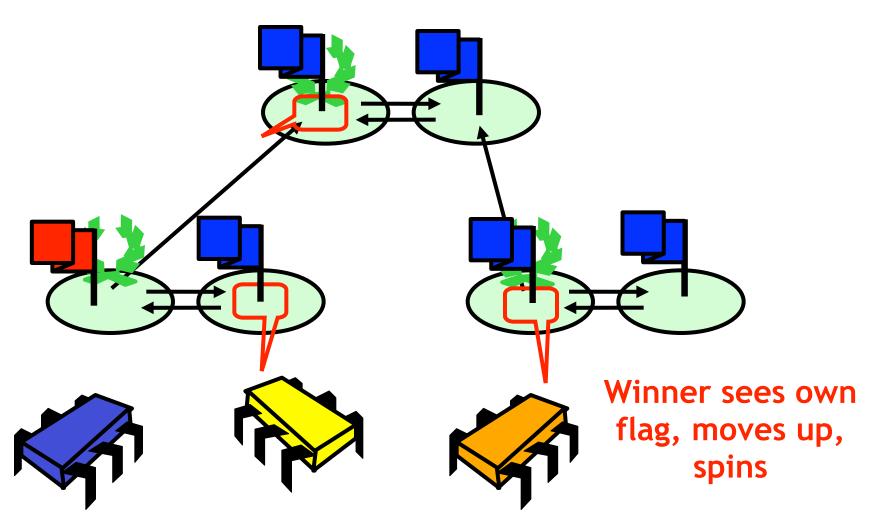


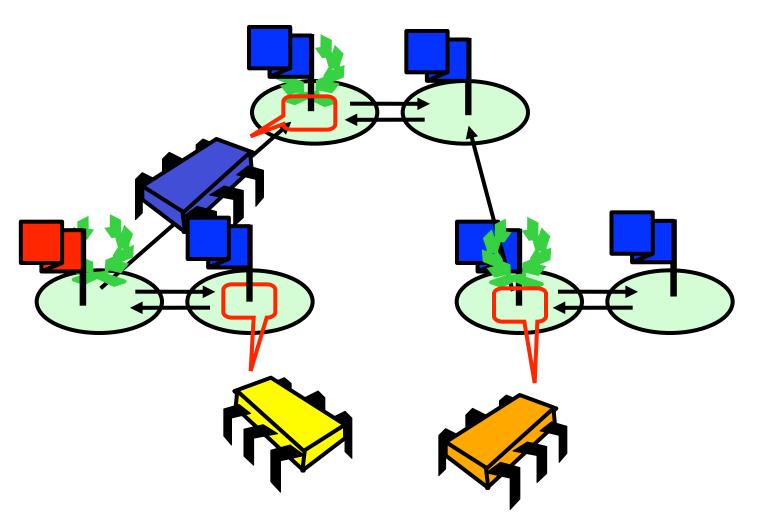


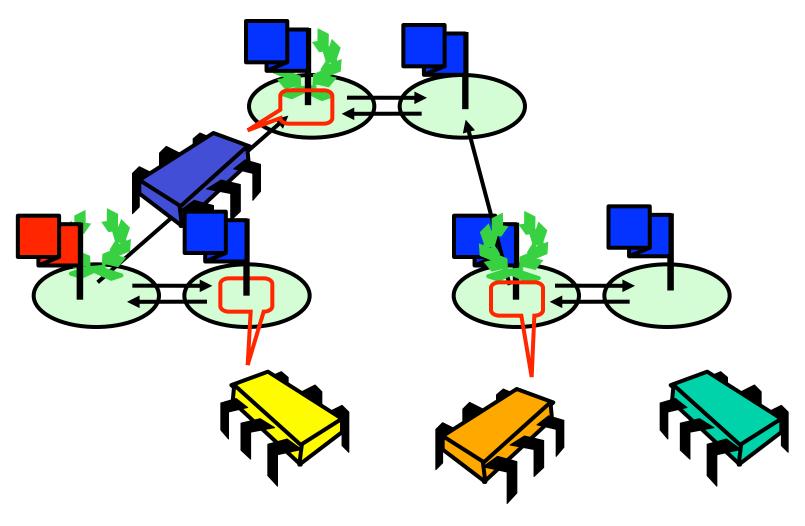


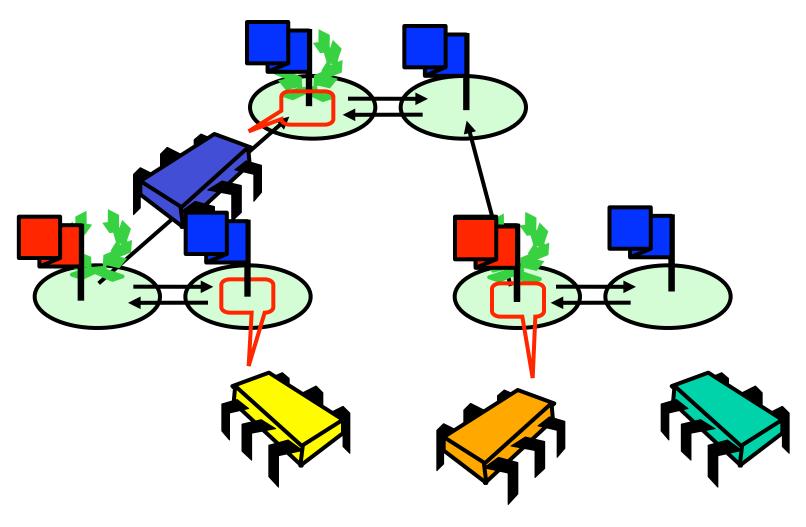


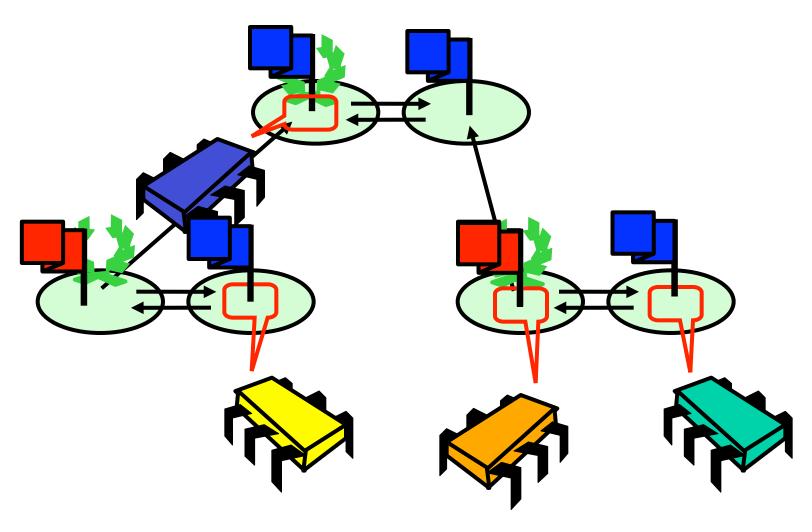


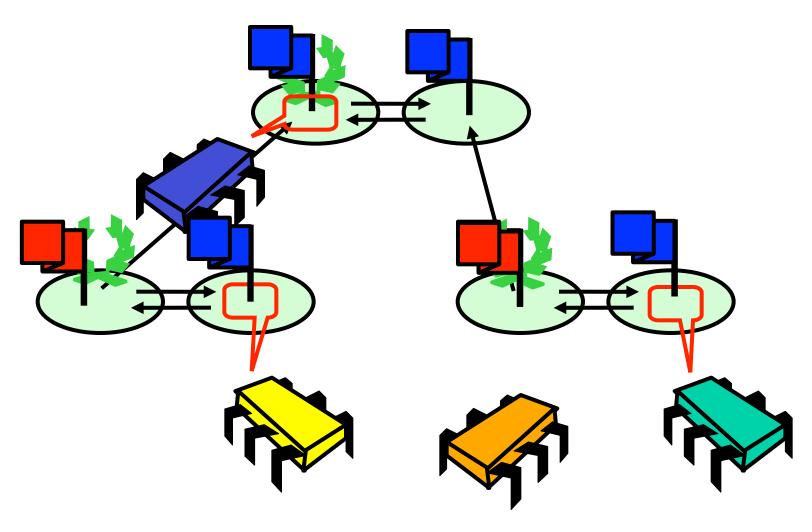


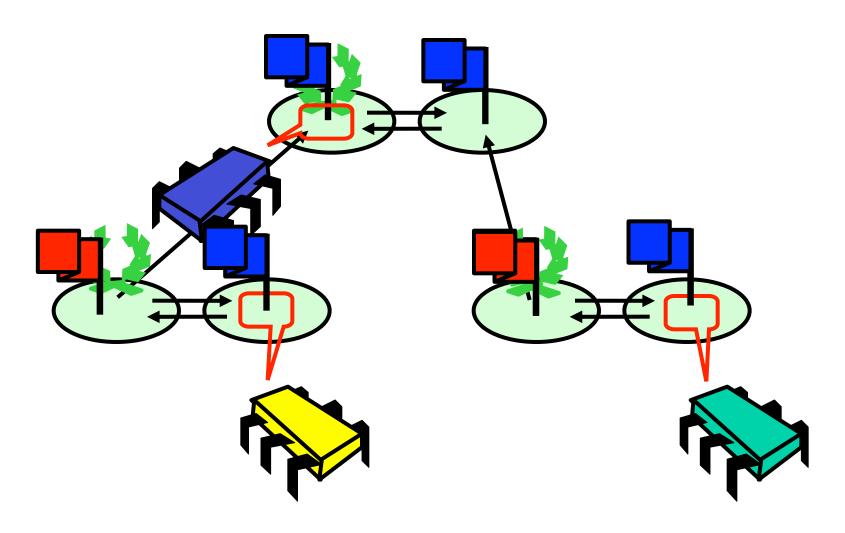


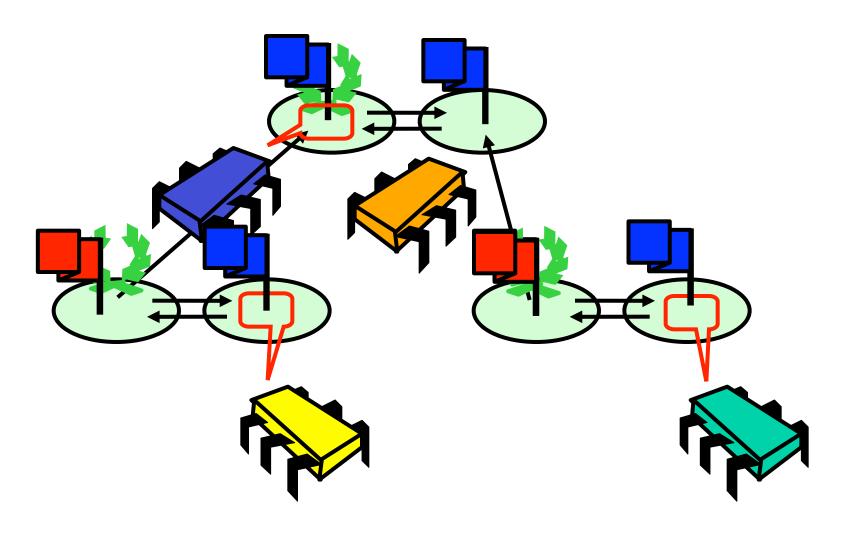


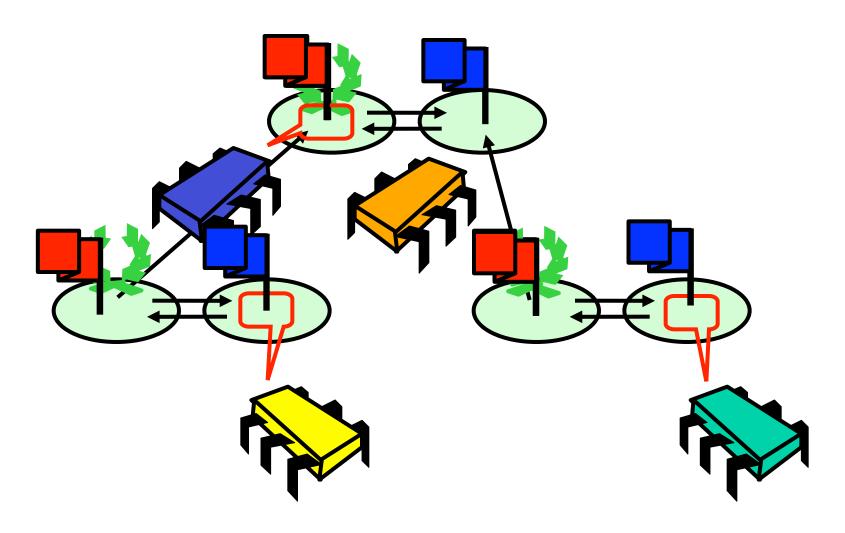


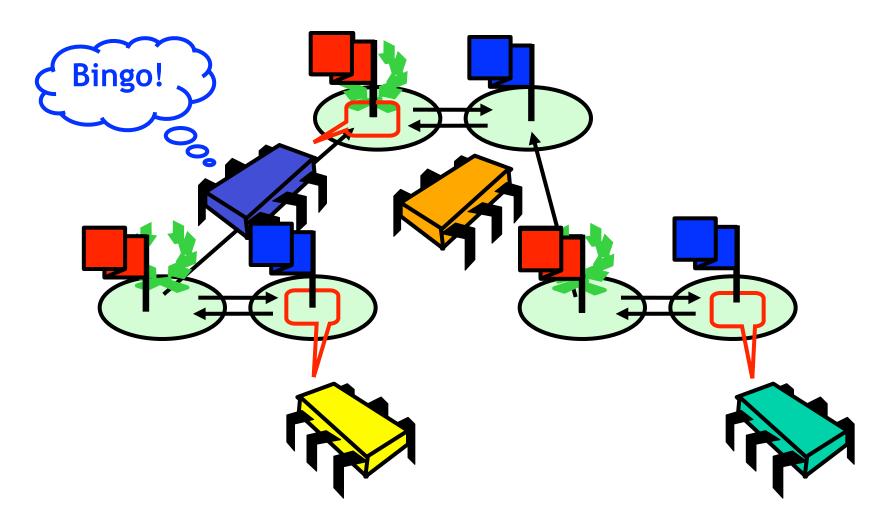


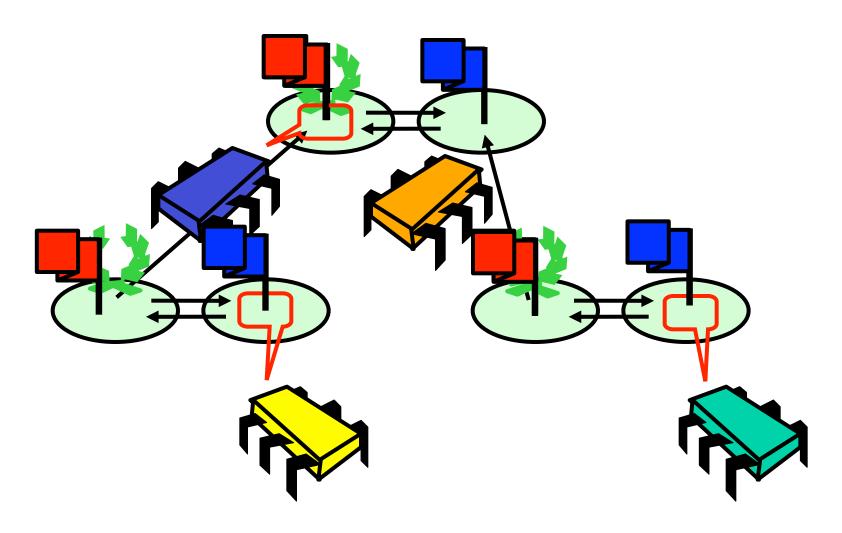


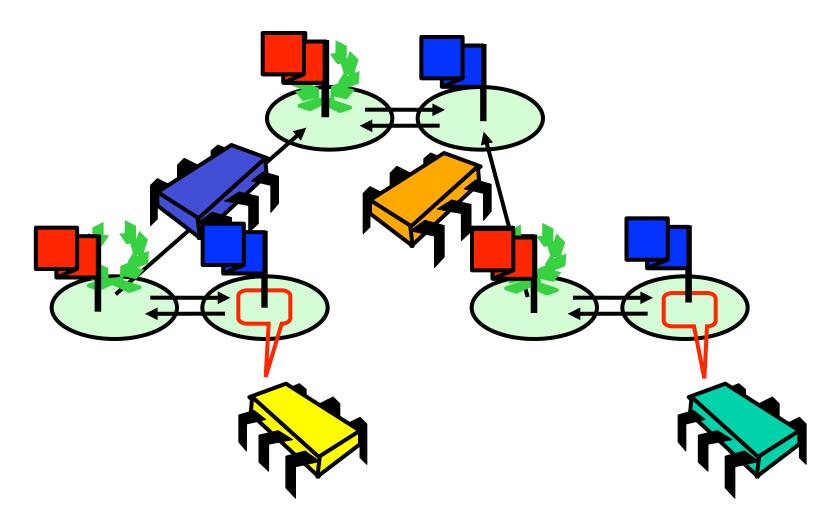


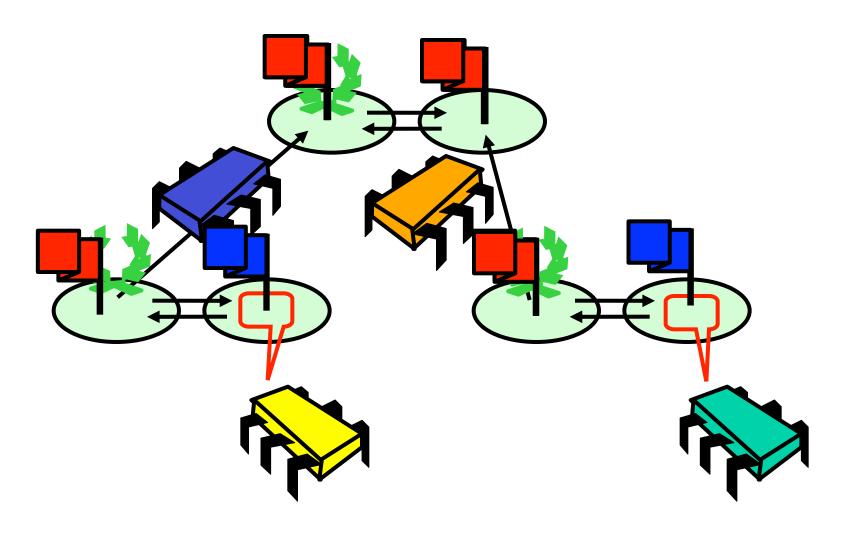


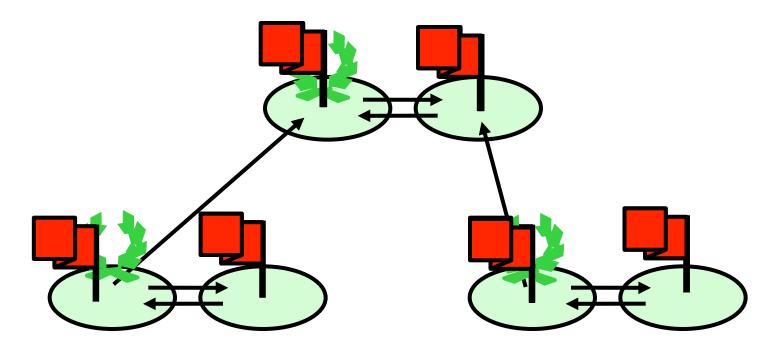












Sense-reversing: next time use blue flags

```
class TBarrier {
  boolean flag;
  TBarrier partner;
  TBarrier parent;
  boolean top;
  ...
}
```

```
boolean flag;
TBarrier partner;
TBarrier parent;
boolean top;
...

Notifications
delivered here
...
```

```
class TBarrier {
   boolean flag;
   TBarrier partner;
   TBarrier parent;
   boolean top;
   ...
}
```

```
class TBarrier {
  boolean flag;
  TBarrier partner;
  TBarrier parent;
  boolean top;
  ...
}
```

```
class TBarrier {
  boolean flag;
  TBarrier parent;
  boolean top;
  ...
}

Am I the root?

Am I the root?
```

```
void await(boolean mySense) {
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySense) {};
 parent.await(mySense);
 partner.flag = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
}}}
```

```
Current sense
void awa t(boolean mySense)
if (top) {
                         Le root, c'est moi
 return;
} else if (parent != null) {
 while (flag != mySense) {};
 parent.await(mySense);
 partner.flag = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                    I am already a winner
if (top) {
 return;
  else if (parent != null) {
 while (flag != mySense) {};
 parent.await(mySense);
 partner.flag = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                          Wait for partner
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySense) {};
 parent.await(mySense);
 partner.flag = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                     Synchronize upstairs
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySense)
 parent.await(mySense);
 partner.flag = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                         Inform partner
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySense) {}
 parent.await(mySeps
partner.flag = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                        Inform partner
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySense) {}
 parent.await(mySeps
 partner.flag = mySense;
} else {
 partner.flag = mySense;
                            Order is important (why?)
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                      Natural-born loser
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySens
 parent.await(mysense);
 partner flo
              = mySense;
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                    Tell partner I'm here
if (top) {
 return;
} else if (parent != null) {
 while (flag != mySense) {};
 parent.await(mySense);
 partner.flag = mySense
} else {
 partner.flag = mySense;
 while (flag != mySense) {};
```

```
void await(boolean mySense) {
                                    Wait for notification
if (top) {
                                         from partner
 return;
} else if (parent != null) {
 while (flag != mySense) {};
 parent.await(mySense);
 partner.flag = mySense;
} else {
 partner.flag = mySense
 while (flag != mySense) {};
```

Remarks

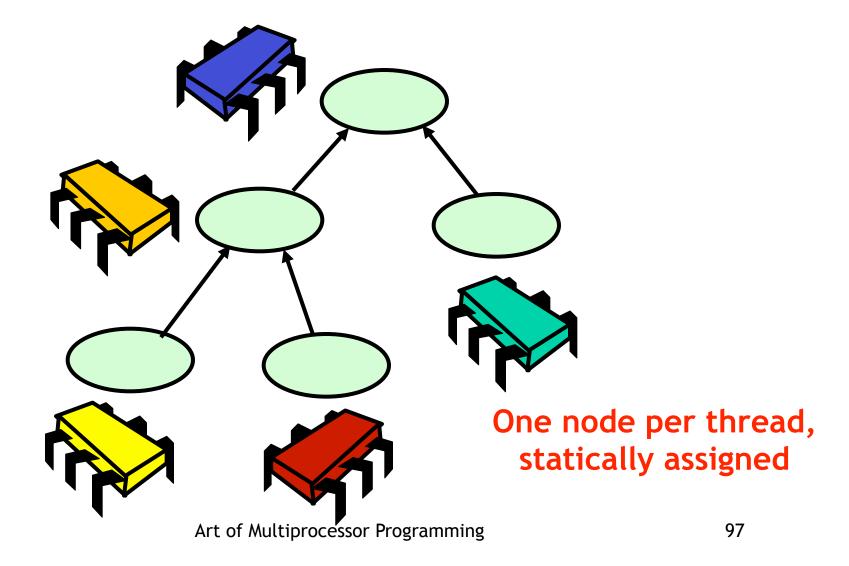
- No need for read-modify-write calls
- Each thread spins on fixed location
 - Good for bus-based architectures
 - Good for NUMA architectures

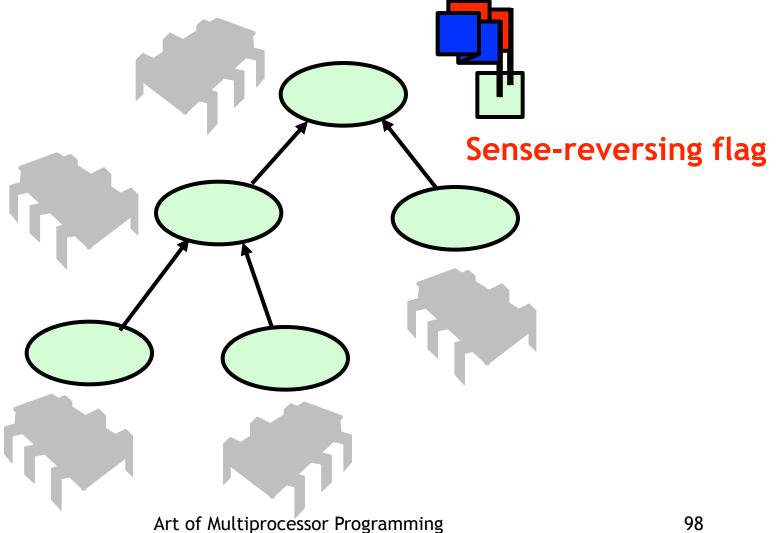
Ideas So Far

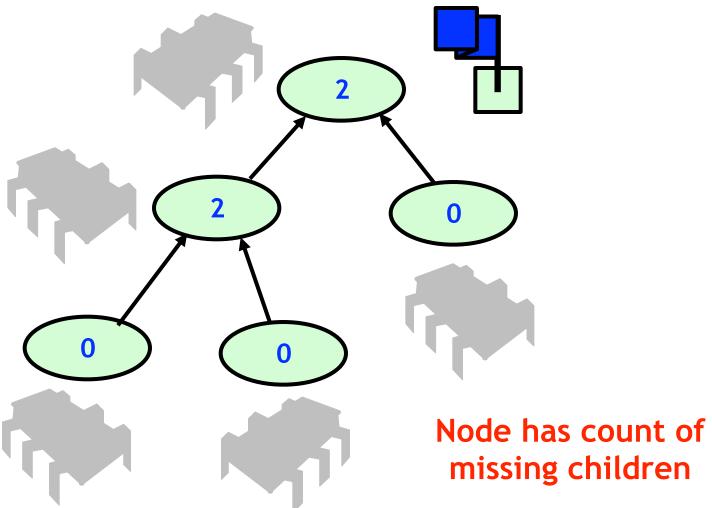
- Sense-reversing
 - Reuse without reinitializing
- Combining tree
 - Like counters, locks ...
- Tournament tree
 - Optimized combining tree

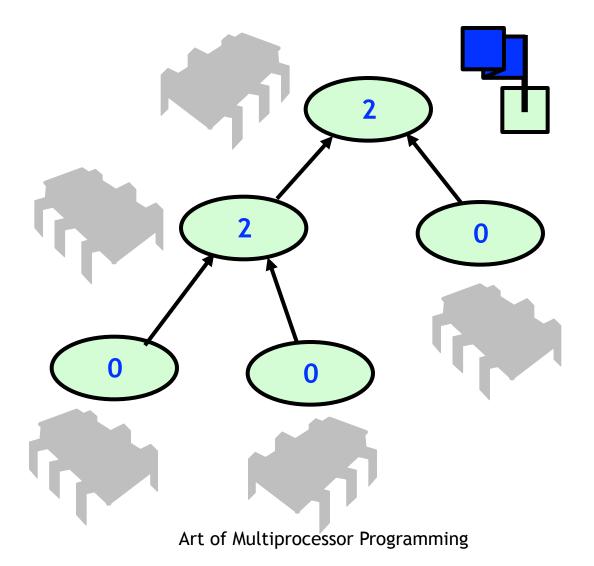
Which is best for Multicore?

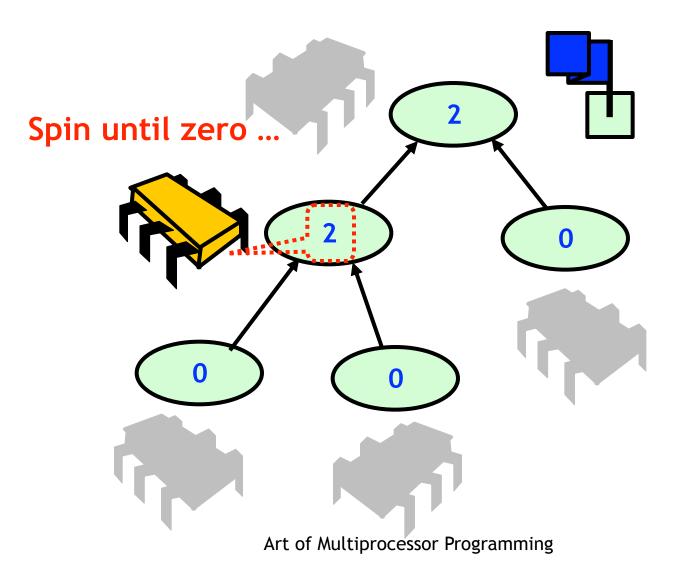
- On a cache coherent multicore chip: perhaps none of the above...
- Here is another (arguably) better algorithm ...

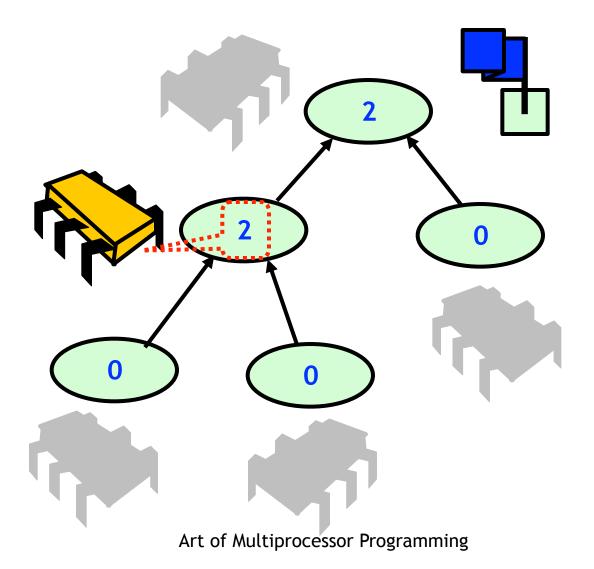


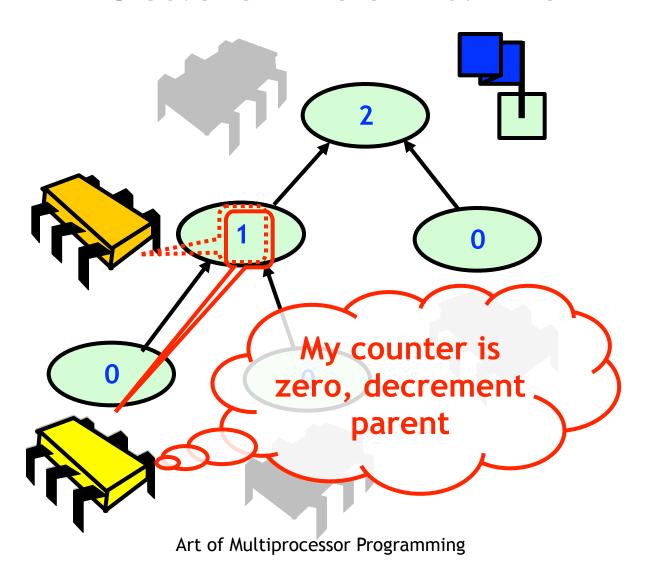


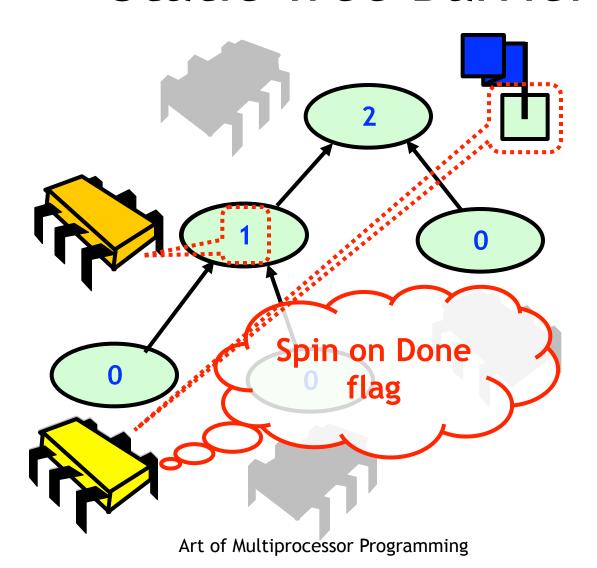


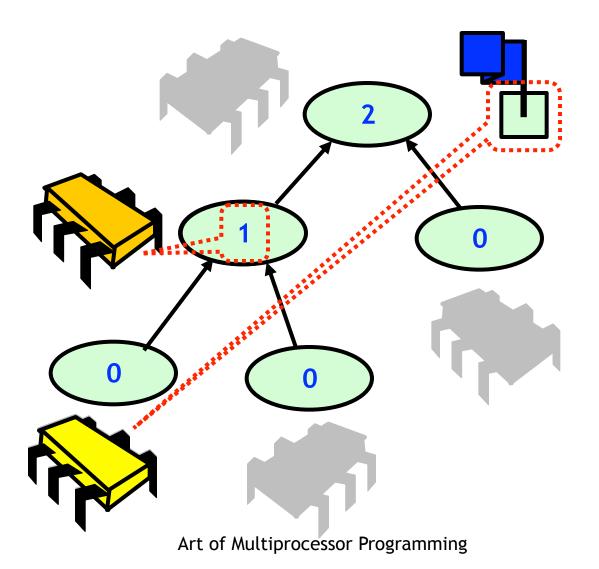


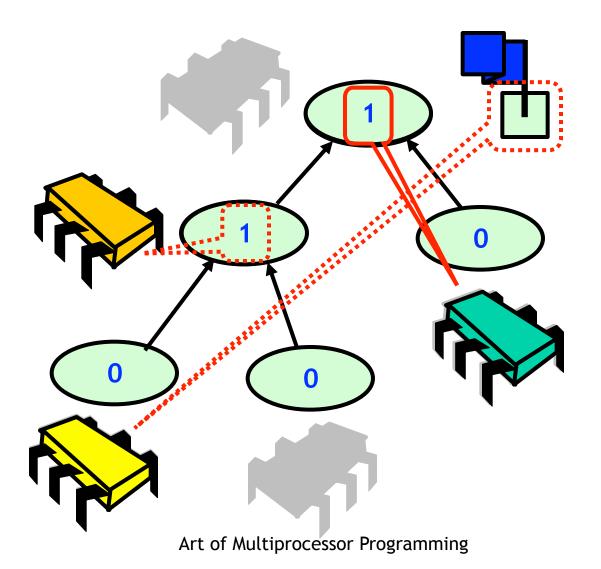


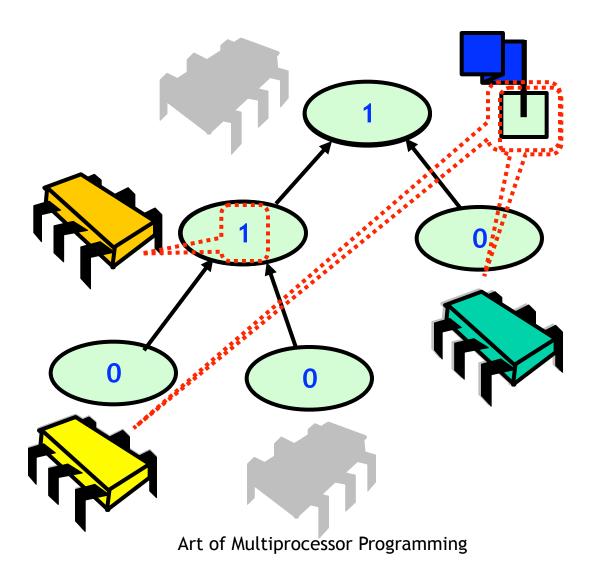


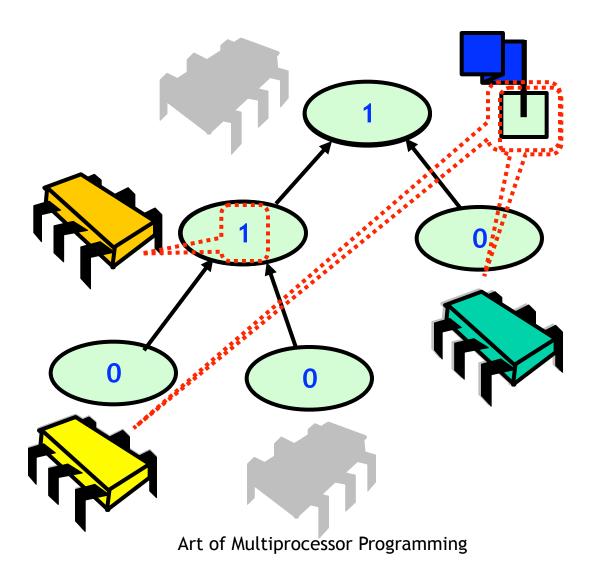


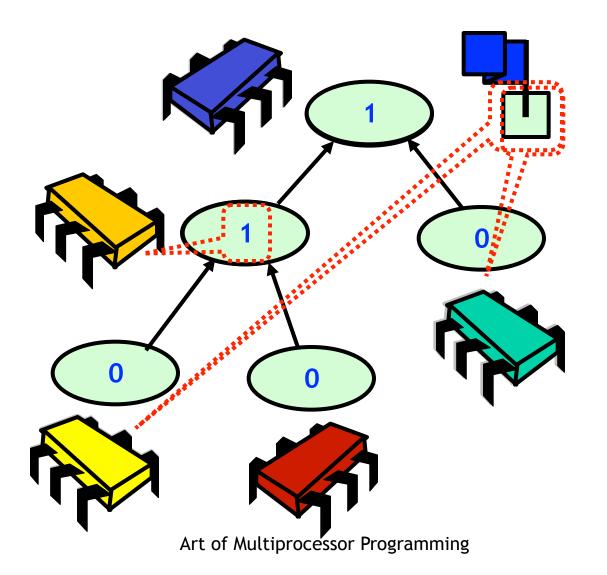


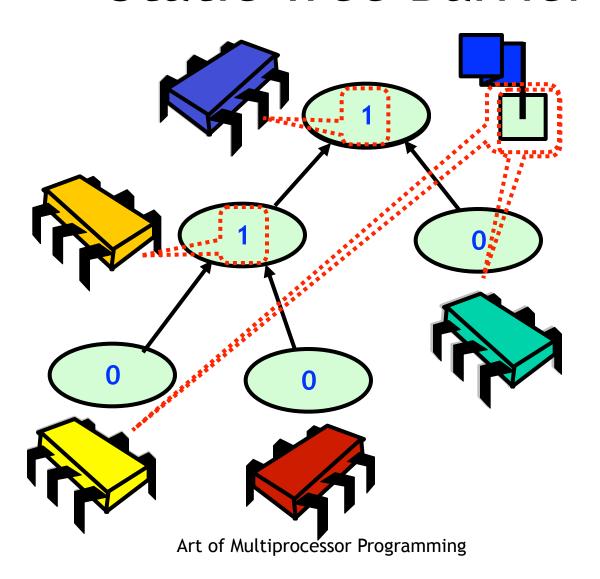


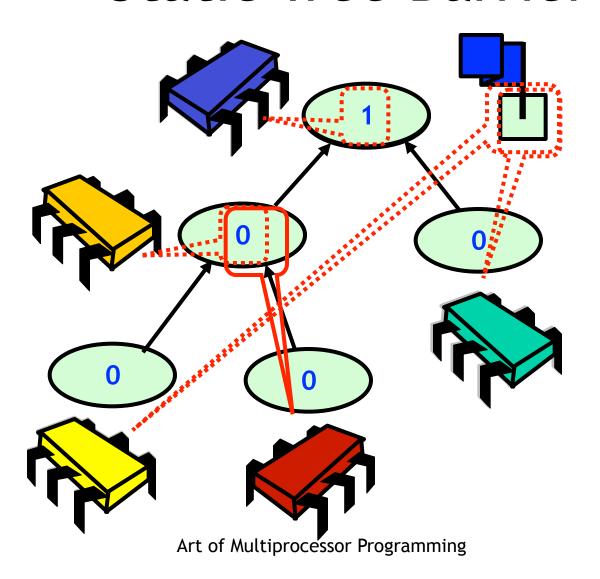


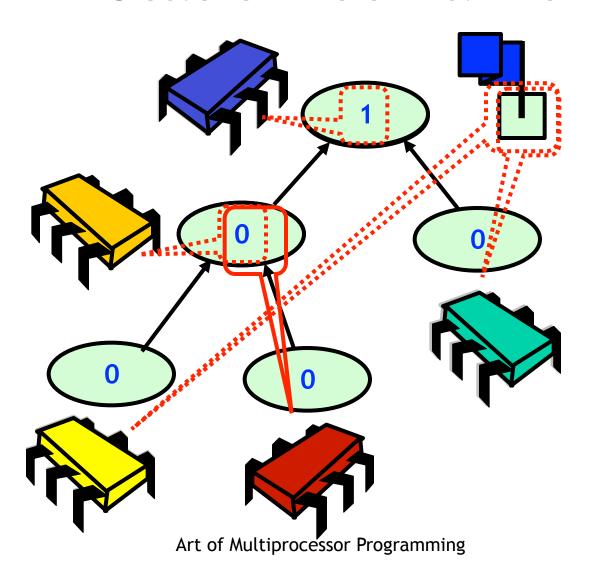


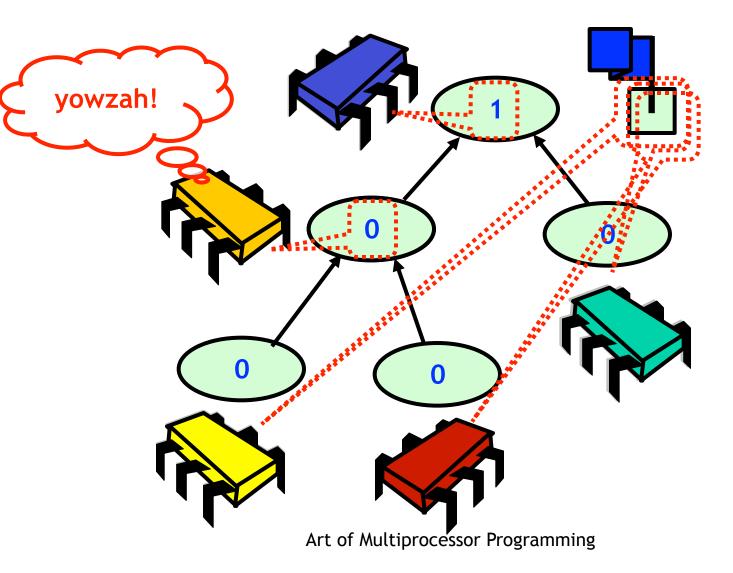


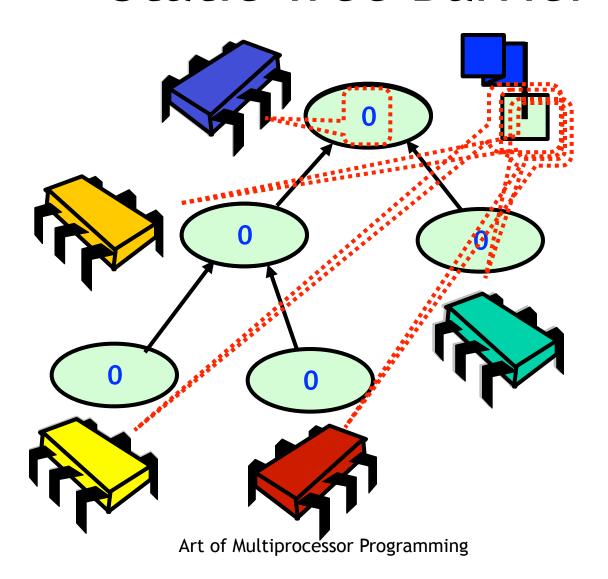






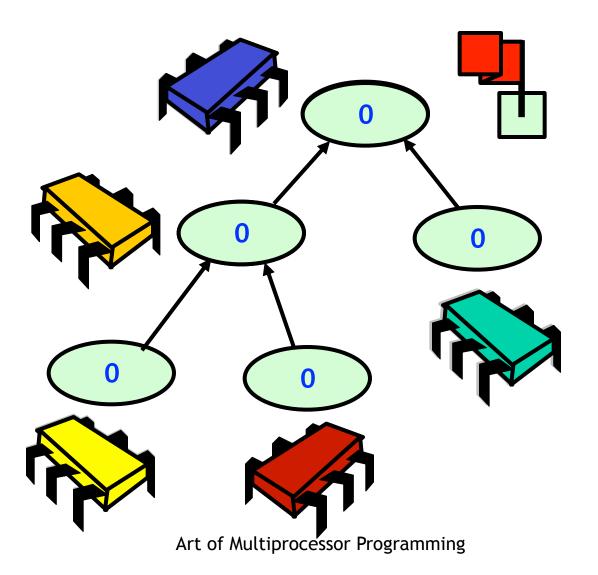




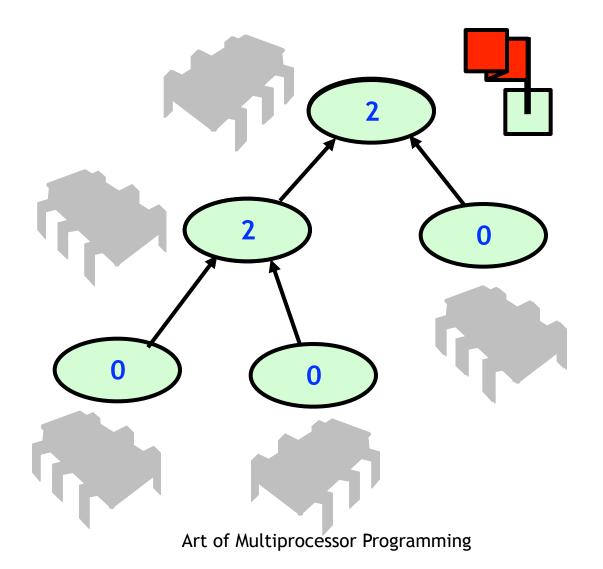


atic Tree Barrier yowzah! 0 Art of Multiprocessor Programming

atic Tree Barrier yowzah! 0 0 Art of Multiprocessor Programming



Static Tree Barrier yes! yes! yes! yes! 0 Art of Multiprocessor Programming 111



Remarks

- Very little cache traffic
- Minimal space overhead
- On message-passing architecture
 - Send notification & sense down tree